

## EURO²

# Basic Quantum Computing Algorithms and Their Implementation in Cirq 

Jiří Tomčala

IT4Innovations,
VŠB - Technical University of Ostrava

5 - 6 September 2023

## Part I

## Introduction to Quantum Computing

## Hardware

Superconducting technology:


## HARDWARE



## Hardware

Trapped-ion technology:


## Hardware



## Qubit

$$
\begin{aligned}
& |\psi\rangle=\alpha|0\rangle+\beta|1\rangle \\
& \alpha=\cos \frac{\theta}{2} \\
& \beta=e^{i \phi} \sin \frac{\theta}{2}=(\cos \phi+i \sin \phi) \sin \frac{\theta}{2} \\
& \operatorname{Pr}(|0\rangle)=|\alpha|^{2}=\cos ^{2} \frac{\theta}{2} \\
& \operatorname{Pr}(|1\rangle)=|\beta|^{2}=\left|e^{i \phi}\right|^{2} \sin ^{2} \frac{\theta}{2}=\sin ^{2} \frac{\theta}{2} \\
& \operatorname{Pr}(|0\rangle)+\operatorname{Pr}(|1\rangle)=\cos ^{2} \frac{\theta}{2}+\sin ^{2} \frac{\theta}{2}=1
\end{aligned}
$$



Figure. Bloch sphere.

## Qubit



Figure. Bloch sphere.

$$
\begin{aligned}
& |\psi\rangle=\alpha|0\rangle+\beta|1\rangle=[|0\rangle|1\rangle]\left[\begin{array}{l}
\alpha \\
\beta
\end{array}\right] \Rightarrow\left[\begin{array}{l}
\alpha \\
\beta
\end{array}\right] \\
& |0\rangle=\left[\begin{array}{l}
1 \\
0
\end{array}\right] \quad|1\rangle=\left[\begin{array}{l}
0 \\
1
\end{array}\right] \\
& X|\psi\rangle=\left[\begin{array}{ll}
0 & 1 \\
1 & 0
\end{array}\right]\left[\begin{array}{l}
\alpha \\
\beta
\end{array}\right]=\left[\begin{array}{l}
\beta \\
\alpha
\end{array}\right] \\
& H|\psi\rangle=\frac{1}{\sqrt{2}}\left[\begin{array}{cc}
1 & 1 \\
1 & -1
\end{array}\right]\left[\begin{array}{l}
\alpha \\
\beta
\end{array}\right]=\frac{1}{\sqrt{2}}\left[\begin{array}{l}
\alpha+\beta \\
\alpha-\beta
\end{array}\right] \\
& H|0\rangle=\frac{1}{\sqrt{2}}\left[\begin{array}{cc}
1 & 1 \\
1 & -1
\end{array}\right]\left[\begin{array}{l}
1 \\
0
\end{array}\right]=\frac{1}{\sqrt{2}}\left[\begin{array}{l}
1 \\
1
\end{array}\right]=|+\rangle \\
& H|1\rangle=\frac{1}{\sqrt{2}}\left[\begin{array}{cc}
1 & 1 \\
1 & -1
\end{array}\right]\left[\begin{array}{l}
0 \\
1
\end{array}\right]=\frac{1}{\sqrt{2}}\left[\begin{array}{c}
1 \\
-1
\end{array}\right]=|-\rangle
\end{aligned}
$$



Figure. Bloch sphere.

## 1-QUBIT QUANTUM GATES

$$
P(\lambda)|\psi\rangle=\left[\begin{array}{cc}
1 & 0 \\
0 & e^{i \lambda}
\end{array}\right]\left[\begin{array}{l}
\alpha \\
\beta
\end{array}\right]=\left[\begin{array}{c}
\alpha \\
e^{i \lambda} \beta
\end{array}\right]
$$

$$
\begin{aligned}
& Z|\psi\rangle=\left[\begin{array}{cc}
1 & 0 \\
0 & e^{i \pi}
\end{array}\right]\left[\begin{array}{l}
\alpha \\
\beta
\end{array}\right]=\left[\begin{array}{cc}
1 & 0 \\
0 & -1
\end{array}\right]\left[\begin{array}{l}
\alpha \\
\beta
\end{array}\right]=\left[\begin{array}{c}
\alpha \\
-\beta
\end{array}\right] \\
& S|\psi\rangle=\left[\begin{array}{cc}
1 & 0 \\
0 & e^{i \frac{\pi}{2}}
\end{array}\right]\left[\begin{array}{l}
\alpha \\
\beta
\end{array}\right]=\left[\begin{array}{cc}
1 & 0 \\
0 & i
\end{array}\right]\left[\begin{array}{l}
\alpha \\
\beta
\end{array}\right]=\left[\begin{array}{c}
\alpha \\
i \beta
\end{array}\right] \\
& T|\psi\rangle=\left[\begin{array}{cc}
1 & 0 \\
0 & e^{i \frac{\pi}{4}}
\end{array}\right]\left[\begin{array}{l}
\alpha \\
\beta
\end{array}\right]=\left[\begin{array}{c}
\alpha \\
\left.e^{i \frac{\pi}{4}} \beta\right]=\left[\begin{array}{c}
\alpha \\
\frac{1}{\sqrt{2}}(1+i) \beta
\end{array}\right]
\end{array} .=\begin{array}{l}
\text { (1) }
\end{array}\right]
\end{aligned}
$$

$$
\begin{aligned}
& Z|+\rangle=|-\rangle \quad Z|-\rangle=|+\rangle \quad S|+\rangle=|i+\rangle \\
& Z|i-\rangle=S|S| i-\rangle=T|T| T|T| i-\rangle=|i+\rangle
\end{aligned}
$$



Figure. Bloch sphere.


## 2-QUBIT QUANTUM GATES

$$
\begin{gathered}
|\psi\rangle=\alpha_{00}|00\rangle+\alpha_{01}|01\rangle+\alpha_{10}|10\rangle+\alpha_{11}|11\rangle=[|00\rangle|01\rangle|10\rangle|11\rangle]\left[\begin{array}{l}
\alpha_{00} \\
\alpha_{01} \\
\alpha_{10} \\
\alpha_{11}
\end{array}\right] \Rightarrow\left[\begin{array}{l}
\alpha_{00} \\
\alpha_{01} \\
\alpha_{10} \\
\alpha_{11}
\end{array}\right] \\
\left|\alpha_{00}\right|^{2}+\left|\alpha_{01}\right|^{2}+\left|\alpha_{10}\right|^{2}+\left|\alpha_{11}\right|^{2}=1 \\
C X|\psi\rangle=\left[\begin{array}{lll}
1 & 0 & 0 \\
0 & 1 & 0 \\
0 & 0 & 0 \\
0 & 0 & 1 \\
0
\end{array}\right]\left[\begin{array}{l}
\alpha_{00} \\
\alpha_{01} \\
\alpha_{10} \\
\alpha_{11}
\end{array}\right]=\left[\begin{array}{l}
\alpha_{00} \\
\alpha_{01} \\
\alpha_{11} \\
\alpha_{10}
\end{array}\right] \\
C P(\lambda)|\psi\rangle=\left[\begin{array}{llll}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 1 & 0 \\
0 & 0 & 0 & e^{i \lambda}
\end{array}\right]\left[\begin{array}{l}
\alpha_{00} \\
\alpha_{01} \\
\alpha_{10} \\
\alpha_{11}
\end{array}\right]=\left[\begin{array}{c}
\alpha_{00} \\
\alpha_{01} \\
\alpha_{10} \\
e^{i \lambda} \alpha_{11}
\end{array}\right] \\
S W A P|\psi\rangle=\left[\begin{array}{lll}
1 & 0 & 0 \\
0 & 0 & 1
\end{array}\right] \\
0 \\
0
\end{gathered} 0
$$



Part II

## Quantum entanglement

$$
\begin{aligned}
& \left.\left.\begin{array}{l}
q_{0}=|0\rangle-H- \\
q_{1}=|0\rangle \longrightarrow-
\end{array}\right\}\left|\psi_{e}\right\rangle=C X|H| 00\right\rangle=C X\left(\frac{1}{\sqrt{2}}|00\rangle+\frac{1}{\sqrt{2}}|01\rangle\right)=\frac{1}{\sqrt{2}}|00\rangle+\frac{1}{\sqrt{2}}|11\rangle=\left|\Phi^{+}\right\rangle
\end{aligned}
$$

$$
\begin{aligned}
& \left.\left.\begin{array}{l}
q_{0}=|0\rangle-H- \\
q_{1}=|0\rangle-\oplus-
\end{array}\right\} \quad\left|\psi_{e}\right\rangle=C X|H| 00\right\rangle=C X\left(\frac{1}{\sqrt{2}}|10\rangle+\frac{1}{\sqrt{2}}|11\rangle\right)=\frac{1}{\sqrt{2}}|10\rangle+\frac{1}{\sqrt{2}}|01\rangle=\left|\Psi^{+}\right\rangle
\end{aligned}
$$



## Part III

## Quantum teleportation



## Implementation in CirQ



## Part IV

## Bernstein-Vazirani + Deutch-Jozsa Algorithm

## BERNSTEIN-VAZIRANI ALGORITHM

The problem statement: Find the secret string $s$ if implemented function f is of the form $f(x)=x \cdot s$.

$$
\begin{aligned}
& |0\rangle^{n} \xrightarrow{H^{\otimes n}} \frac{1}{\sqrt{2^{n}}} \sum_{x \in\{0,1\}^{n}}|x\rangle \xrightarrow{f} \frac{1}{\sqrt{2^{n}}} \sum_{x \in\{0,1\}^{n}}(-1)^{f(x)}|x\rangle \\
& \xrightarrow{H^{\otimes n}} \frac{1}{2^{n}} \sum_{y \in\{0,1\}^{n}} \sum_{x \in\{0,1\}^{n}}(-1)^{f(x)+x \cdot y}|y\rangle=|s\rangle \\
& f(x)+x \cdot y=x \cdot s+x \cdot y=x \cdot(s \oplus y)= \begin{cases}0 & (s=y) \\
0,1,0,1 \ldots & (s \neq y)\end{cases}
\end{aligned}
$$



Figure. Bernstein-Vazirani circuit.


## DEUTCH-JozSA ALGORITHM

The problem statement: Decide whether the implemented function $f$ is constant or balanced.

$$
\begin{aligned}
& |0\rangle^{n} \xrightarrow{H^{\otimes n}} \frac{1}{\sqrt{2^{n}}} \sum_{x \in\{0,1\}^{n}}|x\rangle \xrightarrow{f} \frac{1}{\sqrt{2^{n}}} \sum_{x \in\{0,1\}^{n}}(-1)^{f(x)}|x\rangle \\
& \xrightarrow{H^{\otimes n}} \frac{1}{2^{n}} \sum_{y \in\{0,1\}^{n}} \sum_{x \in\{0,1\}^{n}}(-1)^{f(x)+x \cdot y}|y\rangle=|s\rangle \\
& |s\rangle \begin{cases}=0 \rightarrow f \text { is constant } \\
\neq 0 & \rightarrow f \text { is balanced }\end{cases}
\end{aligned}
$$



Figure. Deutch-Jozsa circuit.


Part V

Simon's ALGORITHM

## Simon's ALGORITHM

The problem statement: Decide whether the implemented function $f$ is periodic or not.

$$
\begin{aligned}
|0\rangle^{\otimes n}|0\rangle^{\otimes n} \xrightarrow{H^{\otimes n}} \frac{1}{\sqrt{2^{n}}} \sum_{x \in\{0,1\}^{n}}|x\rangle|0\rangle^{\otimes n} \\
\xrightarrow{U_{f}} \frac{1}{\sqrt{2^{n}}} \sum_{x \in\{0,1\}^{n}}|x\rangle|f(x)\rangle \\
\xrightarrow{H^{\otimes n}} \frac{1}{2^{n}} \sum_{y \in\{0,1\}^{n}} \sum_{x \in\{0,1\}^{n}}(-1)^{x \cdot y}|y\rangle|f(x)\rangle
\end{aligned}
$$

Quantum state after measuring the lower register:


Figure. Simon's circuit.

$$
f \text { is not periodic } \rightarrow \frac{1}{\sqrt{2^{n}}} \sum_{y \in\{0,1\}^{n}}(-1)^{x_{1} \cdot y}|y\rangle\left|f\left(x_{1}\right)\right\rangle
$$

$f$ is periodic $\rightarrow \frac{1}{\sqrt{2^{n+1+\ldots}}} \sum_{y \in\{0,1\}^{n}}\left[(-1)^{x_{1} \cdot y}+(-1)^{x_{2} \cdot y}+\ldots\right]|y\rangle\left|f\left(x_{1}\right)\right\rangle$

## Implementation in Cirq



## Part VI

Grover's ALGORITHM



## Part VII

## Quantum Fourier transform

## Quantum Fourier transform

$$
\begin{aligned}
& \text { IDFT: } x_{n}=\frac{1}{N} \sum_{k=0}^{N-1} X_{k} \cdot e^{2 \pi i \frac{k n}{N}} \\
& \text { QFT }|x\rangle=\frac{1}{\sqrt{N}} \sum_{y=0}^{N-1} e^{2 \pi i \frac{x y}{N}}|y\rangle \\
& \frac{y}{N}=\frac{y_{1} y_{2} \ldots y_{n}}{2^{n}}=\sum_{k=1}^{n} \frac{y_{k}}{2^{k}} \quad \longrightarrow \quad \mathrm{QFT}|x\rangle=\frac{1}{\sqrt{N}} \sum_{y=0}^{N-1} e^{2 \pi i x \sum_{k=1}^{n} \frac{y_{k}}{2^{k}}}\left|y_{1} y_{2} \ldots y_{n}\right\rangle \\
& \text { QFT }|x\rangle=\frac{1}{\sqrt{2^{n}}} \sum_{y=0}^{2^{n}-1} \prod_{k=1}^{2^{n}} e^{2 \pi i x \frac{y_{k}}{2^{k}}}\left|y_{1} y_{2} \ldots y_{n}\right\rangle \\
& \mathrm{QFT}|x\rangle=\frac{1}{\sqrt{2^{n}}}\left(|0\rangle+e^{i \pi x}|1\rangle\right) \otimes\left(|0\rangle+e^{i \frac{\pi}{2} x}|1\rangle\right) \otimes\left(|0\rangle+e^{i \frac{\pi}{4} x}|1\rangle\right) \otimes \cdots \cdots \otimes\left(|0\rangle+e^{i \frac{\pi}{2^{n-1} x}}|1\rangle\right)
\end{aligned}
$$

Quantum Fourier transform

qubit 3


## Implementation in Cirq

Direct QFT:


## Implementation in Cirq

Inverse QFT:


## Part VIII

## Quantum Phase estimation

## Quantum Phase estimation

## The problem statement:

Estimate the phase of an eigenvalue $e^{2 \pi i \theta}$ of a unitary operator $U$, provided with the corresponding eigenstate $\psi$ :

$$
U|\psi\rangle=e^{2 \pi i \theta}|\psi\rangle
$$


$|0\rangle^{\otimes t} \rightarrow \frac{1}{\sqrt{2^{t}}}(|0\rangle+|1\rangle)^{\otimes t} \rightarrow \frac{1}{\sqrt{2^{t}}}\left(|0\rangle+e^{2 \pi i \theta 2^{t-1}}|1\rangle\right) \otimes\left(|0\rangle+e^{2 \pi i \theta 2^{t-2}}|1\rangle\right) \otimes \cdots \otimes\left(|0\rangle+e^{2 \pi i \theta 2^{0}}|1\rangle\right)=\mathrm{QFT}\left|2^{t} \theta\right\rangle$

## Implementation in Cirq



## Part IX

Shor's ALGORITHM

## Shor's ALGORITHM

## The problem statement:

Find factors $P, R$ of number $N$.

Shor's algorithm procedure:

1. Pick a random integer number $a$ such that: $1<a<N$.
2. If $\operatorname{gcd}(a, N) \neq 1$ then $P=a$ and $R=N / a$.
3. Otherwise, find the period $r$ of function $f(x)=a^{x} \bmod N$.
4. If $r$ is odd then go back to step 1 and choose different $a$.
5. Otherwise, factors $P, R=\operatorname{gcd}\left(a^{r / 2} \pm 1, N\right)$.

A quantum computer can be used for step 3 , in which it is necessary to create a quantum circuit implementing the modular exponentiation function $f(x)=a^{x} \bmod N$ and use this circuit instead of the $U$ operator in the quantum phase estimation circuit.
The resulting circuit is called a period-finder circuit and the measured result at the output can then be used to determine the searched period.

## Shor's ALGORITHM

Period-finder cirquit:


## Implementation in Cirq

Implementation of the function $g(y)=(y \times 6) \bmod 35$ (on the left) and period-finder circuit (on the right) designed to find the period of the function $f(x)=6^{x} \bmod 35$ :


## Thanks

This project has received funding from the European High-Performance Computing Joint Undertaking (JU) under grant agreement No 101101903. The JU receives support from the Digital Europe Programme and Germany, Bulgaria, Austria, Croatia, Cyprus, Czech Republic, Denmark, Estonia, Finland, Greece, Hungary, Ireland, Italy, Lithuania, Latvia, Poland, Portugal, Romania, Slovenia, Spain, Sweden, France, Netherlands, Belgium, Luxembourg, Slovakia, Norway, Türkiye, Republic of North Macedonia, Iceland, Montenegro, Serbia

