



BASIC QUANTUM COMPUTING ALGORITHMS AND THEIR IMPLEMENTATION IN CIRQ

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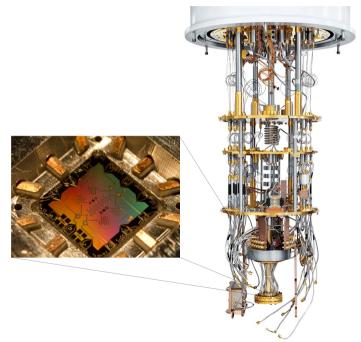
5-6 September 2023

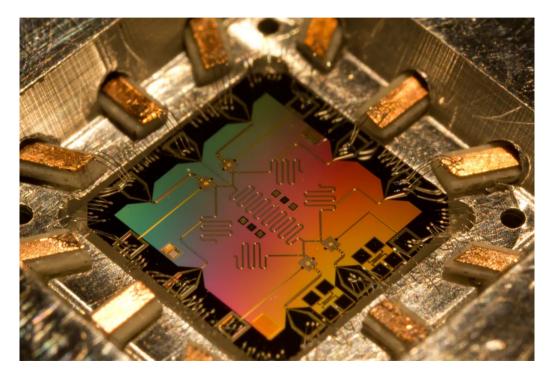


Part I

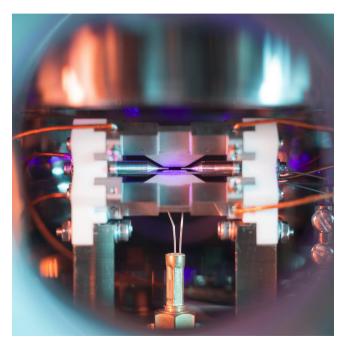
INTRODUCTION TO QUANTUM COMPUTING

Superconducting technology:

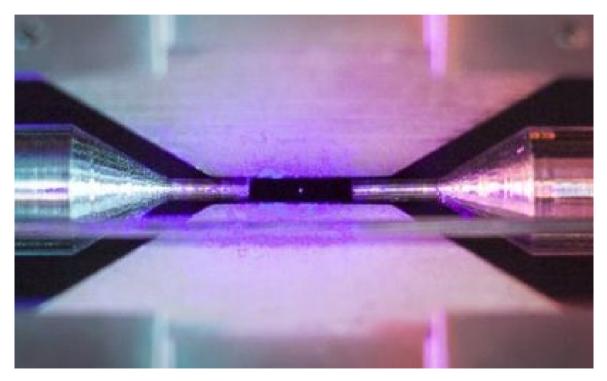




Trapped-ion technology:



BASIC QUANTUM COMPUTING ALGORITHMS AND THEIR IMPLEMENTATION IN CIRQ



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QUBIT

$$\begin{aligned} |\psi\rangle &= \alpha |0\rangle + \beta |1\rangle \\ \alpha &= \cos \frac{\theta}{2} \\ \beta &= e^{i\phi} \sin \frac{\theta}{2} = (\cos \phi + i \sin \phi) \sin \frac{\theta}{2} \\ \Pr(|0\rangle) &= |\alpha|^2 = \cos^2 \frac{\theta}{2} \\ \Pr(|1\rangle) &= |\beta|^2 = |e^{i\phi}|^2 \sin^2 \frac{\theta}{2} = \sin^2 \frac{\theta}{2} \end{aligned}$$

$$\Pr(|0\rangle) + \Pr(|1\rangle) = \cos^2\frac{\theta}{2} + \sin^2\frac{\theta}{2} = 1$$

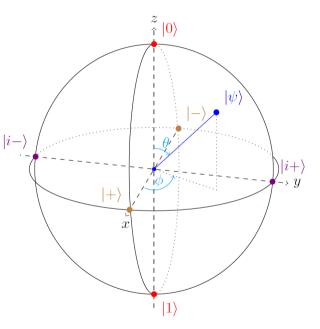


Figure. Bloch sphere.

QUBIT

$$\begin{split} |\psi\rangle &= \alpha \left|0\right\rangle + \beta \left|1\right\rangle \\ \alpha &= \cos\frac{\theta}{2} \\ \beta &= e^{i\phi}\sin\frac{\theta}{2} = (\cos\phi + i\sin\phi)\sin\frac{\theta}{2} \\ |+\rangle &= \frac{1}{\sqrt{2}} \left|0\right\rangle + \frac{1}{\sqrt{2}} \left|1\right\rangle \\ |-\rangle &= \frac{1}{\sqrt{2}} \left|0\right\rangle - \frac{1}{\sqrt{2}} \left|1\right\rangle \\ |i+\rangle &= \frac{1}{\sqrt{2}} \left|0\right\rangle + i\frac{1}{\sqrt{2}} \left|1\right\rangle \\ |i-\rangle &= \frac{1}{\sqrt{2}} \left|0\right\rangle - i\frac{1}{\sqrt{2}} \left|1\right\rangle \end{split}$$

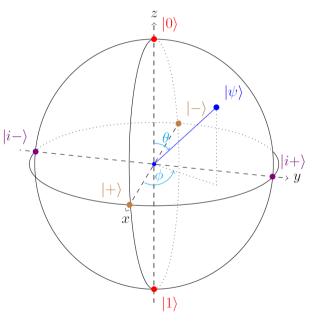


Figure. Bloch sphere.

1-QUBIT QUANTUM GATES

$$\begin{split} |\psi\rangle &= \alpha \left|0\right\rangle + \beta \left|1\right\rangle = \left[|0\rangle \left|1\right\rangle\right] \begin{bmatrix}\alpha\\\beta\end{bmatrix} \Rightarrow \begin{bmatrix}\alpha\\\beta\end{bmatrix}\\\\ |0\rangle &= \begin{bmatrix}1\\0\end{bmatrix} \quad |1\rangle = \begin{bmatrix}0\\1\end{bmatrix}\\\\ X \left|\psi\right\rangle &= \begin{bmatrix}0 & 1\\1 & 0\end{bmatrix} \begin{bmatrix}\alpha\\\beta\end{bmatrix} = \begin{bmatrix}\beta\\\alpha\end{bmatrix}\\\\ H \left|\psi\right\rangle &= \frac{1}{\sqrt{2}} \begin{bmatrix}1 & 1\\1 & -1\end{bmatrix} \begin{bmatrix}\alpha\\\beta\end{bmatrix} = \frac{1}{\sqrt{2}} \begin{bmatrix}\alpha + \beta\\\alpha - \beta\end{bmatrix}\\\\ H \left|0\right\rangle &= \frac{1}{\sqrt{2}} \begin{bmatrix}1 & 1\\1 & -1\end{bmatrix} \begin{bmatrix}1\\0\end{bmatrix} = \frac{1}{\sqrt{2}} \begin{bmatrix}1\\1\end{bmatrix} = |+\rangle\\\\ H \left|1\right\rangle &= \frac{1}{\sqrt{2}} \begin{bmatrix}1 & 1\\1 & -1\end{bmatrix} \begin{bmatrix}0\\1\end{bmatrix} = \frac{1}{\sqrt{2}} \begin{bmatrix}1\\-1\end{bmatrix} = |-\varphi| \end{split}$$

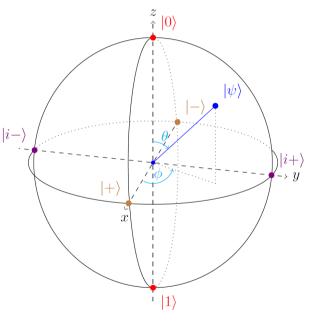


Figure. Bloch sphere.

1-QUBIT QUANTUM GATES

$$P(\lambda) |\psi\rangle = \begin{bmatrix} 1 & 0 \\ 0 & e^{i\lambda} \end{bmatrix} \begin{bmatrix} \alpha \\ \beta \end{bmatrix} = \begin{bmatrix} \alpha \\ e^{i\lambda}\beta \end{bmatrix}$$
$$Z |\psi\rangle = \begin{bmatrix} 1 & 0 \\ 0 & e^{i\pi} \end{bmatrix} \begin{bmatrix} \alpha \\ \beta \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \begin{bmatrix} \alpha \\ \beta \end{bmatrix} = \begin{bmatrix} \alpha \\ -\beta \end{bmatrix}$$
$$S |\psi\rangle = \begin{bmatrix} 1 & 0 \\ 0 & e^{i\frac{\pi}{2}} \end{bmatrix} \begin{bmatrix} \alpha \\ \beta \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \begin{bmatrix} \alpha \\ \beta \end{bmatrix} = \begin{bmatrix} \alpha \\ i\beta \end{bmatrix}$$
$$T |\psi\rangle = \begin{bmatrix} 1 & 0 \\ 0 & e^{i\frac{\pi}{4}} \end{bmatrix} \begin{bmatrix} \alpha \\ \beta \end{bmatrix} = \begin{bmatrix} \alpha \\ e^{i\frac{\pi}{4}}\beta \end{bmatrix} = \begin{bmatrix} \alpha \\ \frac{1}{\sqrt{2}}(1+i)\beta \end{bmatrix}$$
$$Z |+\rangle = |-\rangle \quad Z |-\rangle = |+\rangle \quad S |+\rangle = |i+\rangle$$
$$Z |i-\rangle = S|S |i-\rangle = T|T|T|T |i-\rangle = |i+\rangle$$

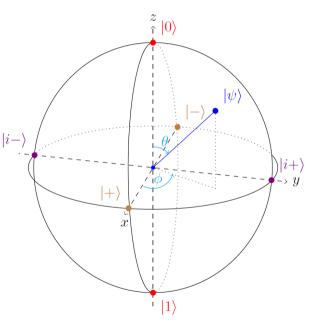
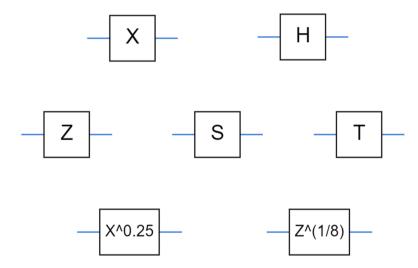


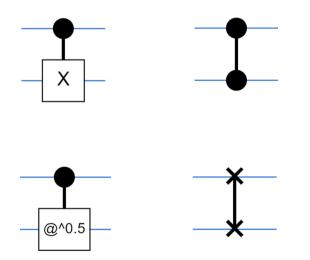
Figure. Bloch sphere.



2-QUBIT QUANTUM GATES

$$\begin{split} \psi \rangle &= \alpha_{00} |00\rangle + \alpha_{01} |01\rangle + \alpha_{10} |10\rangle + \alpha_{11} |11\rangle = \left[|00\rangle |01\rangle |10\rangle |11\rangle \right] \begin{bmatrix} \alpha_{00} \\ \alpha_{01} \\ \alpha_{10} \\ \alpha_{11} \end{bmatrix} \Rightarrow \begin{bmatrix} \alpha_{00} \\ \alpha_{01} \\ \alpha_{10} \\ \alpha_{11} \end{bmatrix} \\ |\alpha_{00}|^2 + |\alpha_{01}|^2 + |\alpha_{10}|^2 + |\alpha_{11}|^2 = 1 \\ CX |\psi\rangle &= \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & e^{i\lambda} \end{bmatrix} \begin{bmatrix} \alpha_{00} \\ \alpha_{01} \\ \alpha_{11} \\ \alpha_{10} \end{bmatrix} = \begin{bmatrix} \alpha_{00} \\ \alpha_{01} \\ \alpha_{11} \\ \alpha_{10} \end{bmatrix} \\ CP(\lambda) |\psi\rangle &= \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & e^{i\lambda} \end{bmatrix} \begin{bmatrix} \alpha_{00} \\ \alpha_{01} \\ \alpha_{11} \\ \alpha_{10} \\ \alpha_{11} \end{bmatrix} = \begin{bmatrix} \alpha_{00} \\ \alpha_{01} \\ \alpha_{10} \\ \alpha_{10} \\ \alpha_{11} \end{bmatrix} \\ SWAP |\psi\rangle &= \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} \alpha_{00} \\ \alpha_{01} \\ \alpha_{10} \\ \alpha_{11} \\ \alpha_{11} \end{bmatrix} = \begin{bmatrix} \alpha_{00} \\ \alpha_{01} \\ \alpha_{10} \\ \alpha_{11} \\ \alpha_{11} \end{bmatrix} \end{split}$$

BASIC QUANTUM COMPUTING ALGORITHMS AND THEIR IMPLEMENTATION IN CIRQ





Part II

QUANTUM ENTANGLEMENT

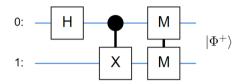
Bell states

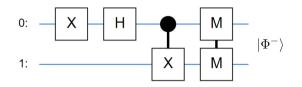
$$\begin{array}{c} q_0 = |0\rangle \quad -\underline{H} \\ \hline \bullet \\ q_1 = |0\rangle \quad -\underline{\bullet} \end{array} \right\} \quad |\psi_e\rangle = CX \left| H \left| 00 \right\rangle = CX \left(\frac{1}{\sqrt{2}} \left| 00 \right\rangle + \frac{1}{\sqrt{2}} \left| 01 \right\rangle \right) = \frac{1}{\sqrt{2}} \left| 00 \right\rangle + \frac{1}{\sqrt{2}} \left| 11 \right\rangle = |\Phi^+\rangle$$

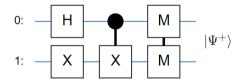
$$\begin{array}{c} q_0 = |0\rangle & \textcircled{H} \\ q_1 = |0\rangle & \textcircled{H} \\ \end{array} \right\} \quad |\psi_e\rangle = CX|H|00\rangle = CX\left(\frac{1}{\sqrt{2}}|00\rangle - \frac{1}{\sqrt{2}}|01\rangle\right) = \frac{1}{\sqrt{2}}|00\rangle - \frac{1}{\sqrt{2}}|11\rangle = |\Phi^-\rangle$$

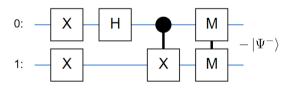
$$\begin{array}{c} q_0 = |0\rangle \quad -\underline{H} \\ \hline \bullet \\ q_1 = |0\rangle \quad -\underline{\bullet} \\ \end{array} \right\} \quad |\psi_e\rangle = CX|H|00\rangle = CX\left(\frac{1}{\sqrt{2}}|10\rangle + \frac{1}{\sqrt{2}}|11\rangle\right) = \frac{1}{\sqrt{2}}|10\rangle + \frac{1}{\sqrt{2}}|01\rangle = |\Psi^+\rangle$$

$$\begin{array}{c} q_0 = |0\rangle & \textcircled{H} & \textcircled{H} \\ q_1 = |0\rangle & \textcircled{H} & \textcircled{H} \end{array} \end{array} \right\} \quad |\psi_e\rangle = CX|H|00\rangle = CX\left(\frac{1}{\sqrt{2}}|10\rangle - \frac{1}{\sqrt{2}}|11\rangle\right) = \frac{1}{\sqrt{2}}|10\rangle - \frac{1}{\sqrt{2}}|01\rangle = -|\Psi^-\rangle$$





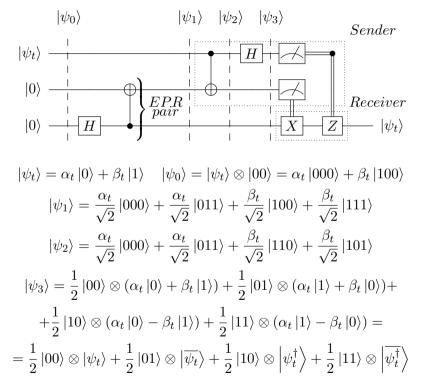




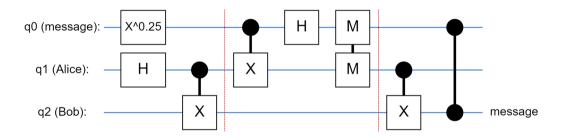


Part III

QUANTUM TELEPORTATION



BASIC QUANTUM COMPUTING ALGORITHMS AND THEIR IMPLEMENTATION IN CIRQ





Part IV

${\tt Bernstein-Vazirani} + {\tt Deutch-Jozsa \ algorithm}$

BERNSTEIN-VAZIRANI ALGORITHM

The problem statement: Find the secret string s if implemented function f is of the form $f(x) = x \cdot s$.

$$\begin{aligned} |0\rangle^n \xrightarrow{H^{\otimes n}} \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} |x\rangle & \xrightarrow{f} \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} (-1)^{f(x)} |x\rangle \\ & \xrightarrow{H^{\otimes n}} \frac{1}{2^n} \sum_{y \in \{0,1\}^n} \sum_{x \in \{0,1\}^n} (-1)^{f(x)+x \cdot y} |y\rangle = |s\rangle \end{aligned}$$

$$\begin{aligned} f'(x) + x \cdot y = x \cdot s + x \cdot y = x \cdot (s \oplus y) = \begin{cases} 0 & (s = y) \\ 0, 1, 0, 1 \dots & (s \neq y) \end{cases} \end{aligned}$$

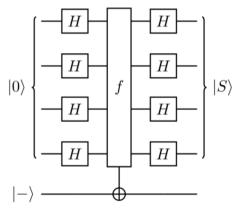
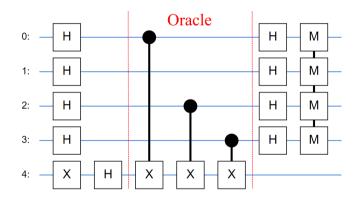


Figure. Bernstein-Vazirani circuit.



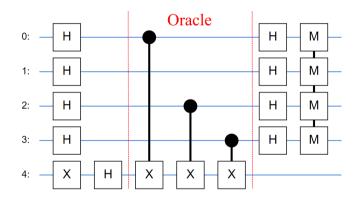
DEUTCH-JOZSA ALGORITHM

The problem statement: Decide whether the implemented function f is constant or balanced.

$$\begin{split} |0\rangle^n \xrightarrow{H^{\otimes n}} \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} |x\rangle \xrightarrow{f} \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} (-1)^{f(x)} |x\rangle \\ \xrightarrow{H^{\otimes n}} \frac{1}{2^n} \sum_{y \in \{0,1\}^n} \sum_{x \in \{0,1\}^n} (-1)^{f(x)+x \cdot y} |y\rangle = |s\rangle \\ |s\rangle \begin{cases} = 0 \ \to f \text{ is constant} \\ \neq 0 \ \to f \text{ is balanced} \end{cases} \end{split}$$

Figure. Deutch-Jozsa circuit.

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Part V

SIMON'S ALGORITHM

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SIMON'S ALGORITHM

The problem statement: Decide whether the implemented function f is periodic or not.

$$\begin{aligned} |0\rangle^{\otimes n} |0\rangle^{\otimes n} \xrightarrow{H^{\otimes n}} \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} |x\rangle |0\rangle^{\otimes n} \\ \xrightarrow{U_f} \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} |x\rangle |f(x)\rangle \\ \xrightarrow{H^{\otimes n}} \frac{1}{2^n} \sum_{y \in \{0,1\}^n} \sum_{x \in \{0,1\}^n} (-1)^{x \cdot y} |y\rangle |f(x)\rangle \end{aligned}$$

Quantum state after measuring the lower register:

$$f \text{ is not periodic } \rightarrow \frac{1}{\sqrt{2^n}} \sum_{y \in \{0,1\}^n} (-1)^{x_1 \cdot y} \ket{y} \ket{f(x_1)}$$

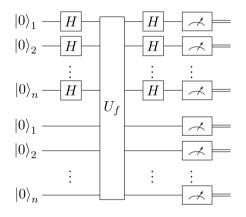
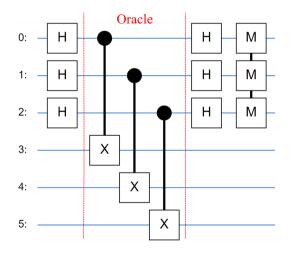


Figure. Simon's circuit.

$$f \text{ is periodic } \rightarrow \frac{1}{\sqrt{2^{n+1+\dots}}} \sum_{y \in \{0,1\}^n} \left[(-1)^{x_1 \cdot y} + (-1)^{x_2 \cdot y} + \dots \right] |y\rangle |f(x_1)\rangle$$

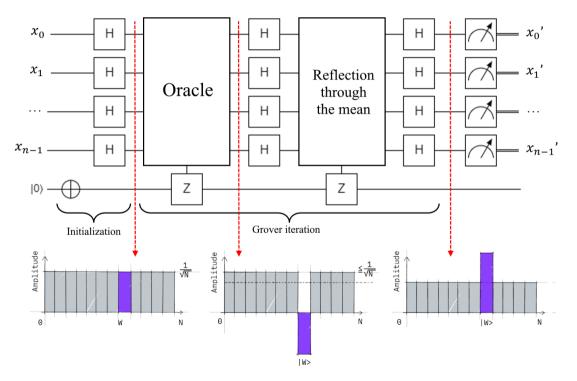


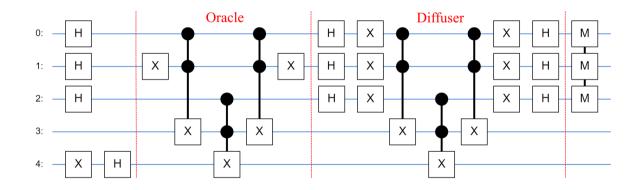


Part VI

GROVER'S ALGORITHM

GROVER'S ALGORITHM







Part VII

QUANTUM FOURIER TRANSFORM

QUANTUM FOURIER TRANSFORM

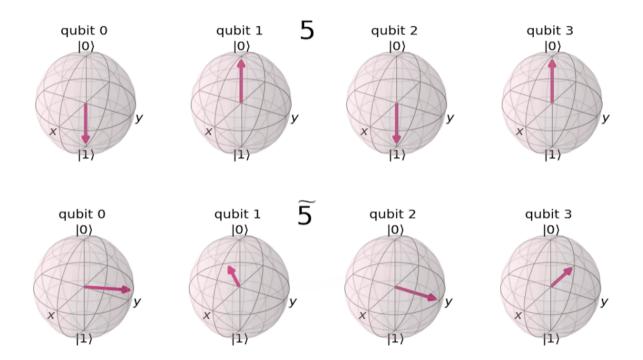
IDFT:
$$x_n = \frac{1}{N} \sum_{k=0}^{N-1} X_k \cdot e^{2\pi i \frac{kn}{N}}$$

$$\operatorname{QFT} \left| x \right\rangle = \frac{1}{\sqrt{N}} \sum_{y=0}^{N-1} e^{2\pi i \frac{xy}{N}} \left| y \right\rangle$$

$$\frac{y}{N} = \frac{y_1 y_2 \dots y_n}{2^n} = \sum_{k=1}^n \frac{y_k}{2^k} \quad \longrightarrow \quad \text{QFT} |x\rangle = \frac{1}{\sqrt{N}} \sum_{y=0}^{N-1} e^{2\pi i x \sum_{k=1}^n \frac{y_k}{2^k}} |y_1 y_2 \dots y_n\rangle$$

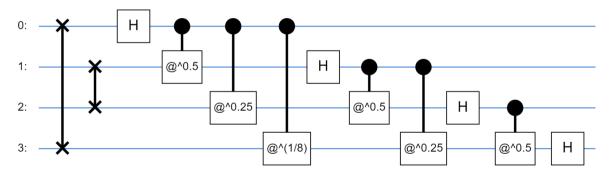
$$\operatorname{QFT}|x\rangle = \frac{1}{\sqrt{2^{n}}} \Big(|0\rangle + e^{i\pi x} |1\rangle\Big) \otimes \Big(|0\rangle + e^{i\frac{\pi}{2}x} |1\rangle\Big) \otimes \Big(|0\rangle + e^{i\frac{\pi}{4}x} |1\rangle\Big) \otimes \cdots \otimes \Big(|0\rangle + e^{i\frac{\pi}{2^{n-1}}x} |1\rangle\Big)$$

QUANTUM FOURIER TRANSFORM

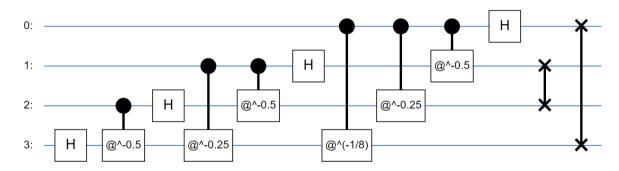


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Direct QFT:



Inverse QFT:





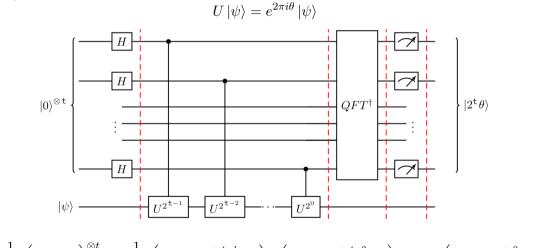
Part VIII

QUANTUM PHASE ESTIMATION

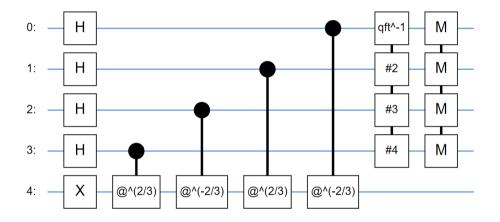
QUANTUM PHASE ESTIMATION

The problem statement:

Estimate the phase of an eigenvalue $e^{2\pi i\theta}$ of a unitary operator U, provided with the corresponding eigenstate ψ :



$$|0\rangle^{\otimes t} \rightarrow \frac{1}{\sqrt{2^{t}}} \Big(|0\rangle + |1\rangle\Big)^{\otimes t} \rightarrow \frac{1}{\sqrt{2^{t}}} \Big(|0\rangle + e^{2\pi i \theta 2^{t-1}} |1\rangle\Big) \otimes \Big(|0\rangle + e^{2\pi i \theta 2^{t-2}} |1\rangle\Big) \otimes \cdots \otimes \Big(|0\rangle + e^{2\pi i \theta 2^{0}} |1\rangle\Big) = \operatorname{QFT} |2^{t}\theta\rangle$$





Part IX

SHOR'S ALGORITHM

SHOR'S ALGORITHM

The problem statement:

Find factors P, R of number N.

Shor's algorithm procedure:

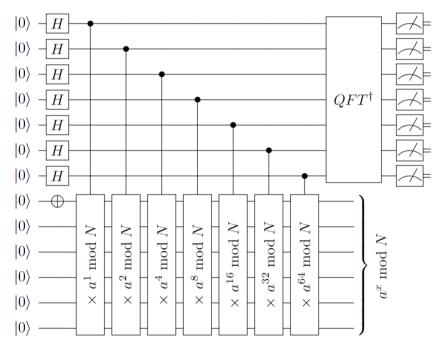
- 1. Pick a random integer number a such that: 1 < a < N.
- 2. If $gcd(a, N) \neq 1$ then P = a and R = N/a.
- 3. Otherwise, find the period r of function $f(x) = a^x \mod N$.
- 4. If r is odd then go back to step 1 and choose different a.
- 5. Otherwise, factors $P, R = \text{gcd}(a^{r/2} \pm 1, N)$.

A quantum computer can be used for step 3, in which it is necessary to create a quantum circuit implementing the modular exponentiation function $f(x) = a^x \mod N$ and use this circuit instead of the U operator in the quantum phase estimation circuit.

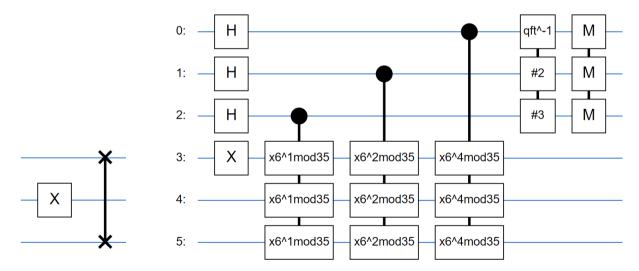
The resulting circuit is called a period-finder circuit and the measured result at the output can then be used to determine the searched period.

SHOR'S ALGORITHM

Period-finder cirquit:



Implementation of the function $g(y) = (y \times 6) \mod 35$ (on the left) and period-finder circuit (on the right) designed to find the period of the function $f(x) = 6^x \mod 35$:





Thanks



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