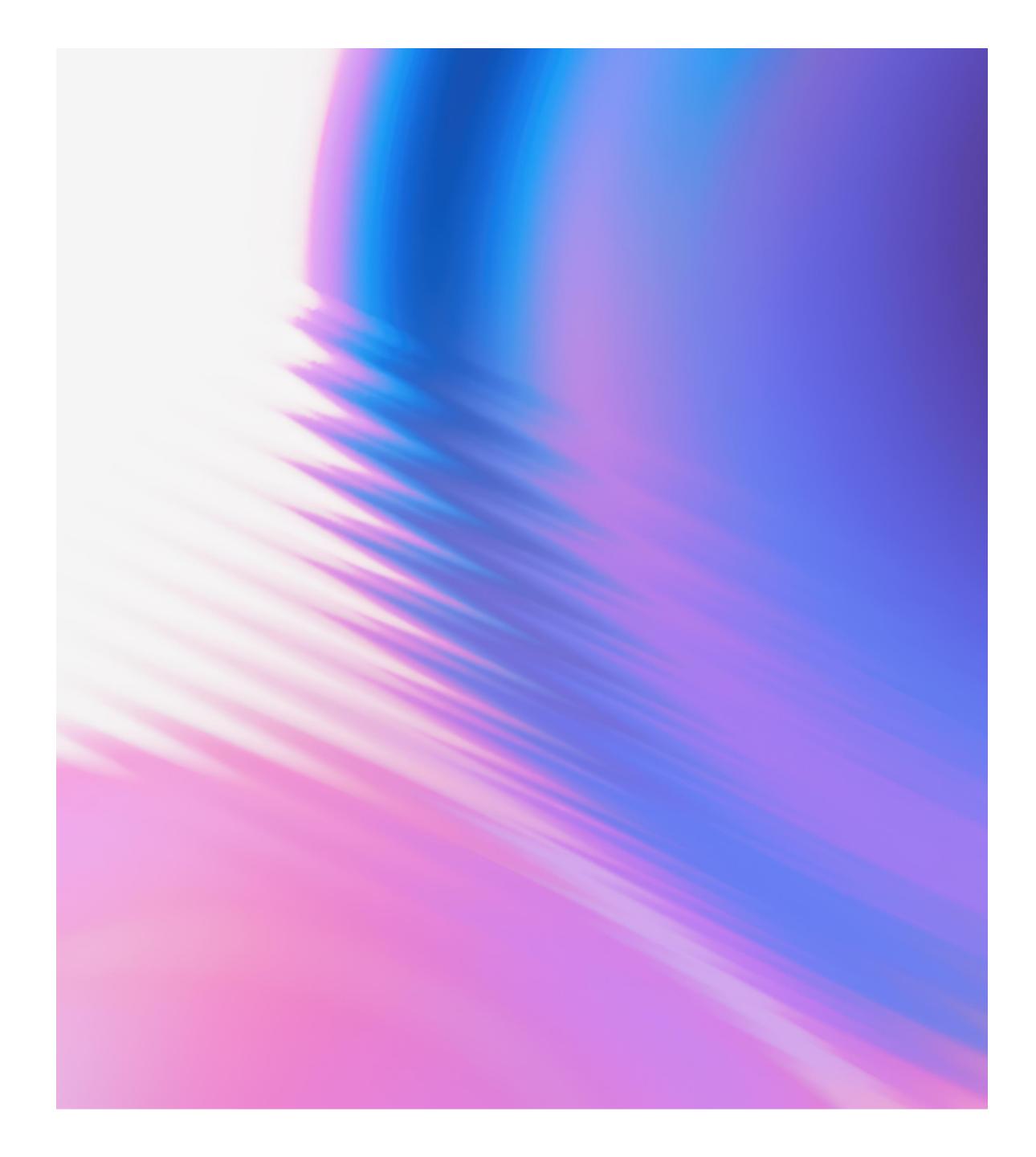
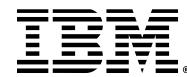
# Dynamic Circuits for Efficient Quantum Computation

Elisa Bäumer Marty Research Scientist IBM Quantum

#### Outline:

- IBM Quantum Roadmap
- Dynamic Circuits
- Long-Range Entanglement
- Quantum Fourier Transform
- > Logarithmic-Depth AQFT on a Line





Key milestones in bringing useful quantum computing to the world

2023

Establish quantum utility



2026

Demonstrate quantum advantage



2029

Deliver the first largescale, fault-tolerant quantum computer







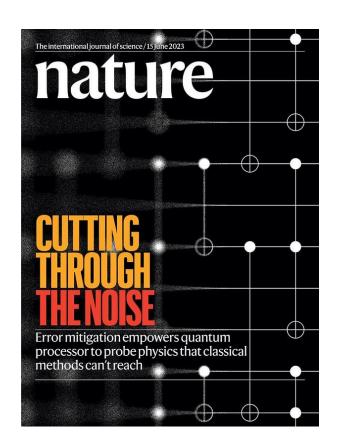


# Quantum Utility (2023)



Demonstration that a quantum computer can run quantum circuits beyond the ability of a classical computer simulating a quantum computer

Confirmation via research, papers, & theory



IBM's 2023 research paper ("Evidence for the utility of quantum computing before fault tolerance") provided evidence and methods to move the industry into the Utility era

https://www.nature.com/articles/s41586-023-06096-3

### Quantum Advantage (TBD)



Demonstration that a quantum computer can solve a problem more accurately, cheaper, or more efficiently than classical computing alone

#### Confirmation via real-world usage



Advantage will come at different times in different domains and depends on the continued advancement of quantum algorithm implementations across industries

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Performing an information processing task more efficiently, cost-effectively, or accurately using quantum computers than is known to be possible with classical computers alone

1

We must establish trust in the outputs of a real and noisy quantum device.

A scientific experiment is only as good as confidence in methods.

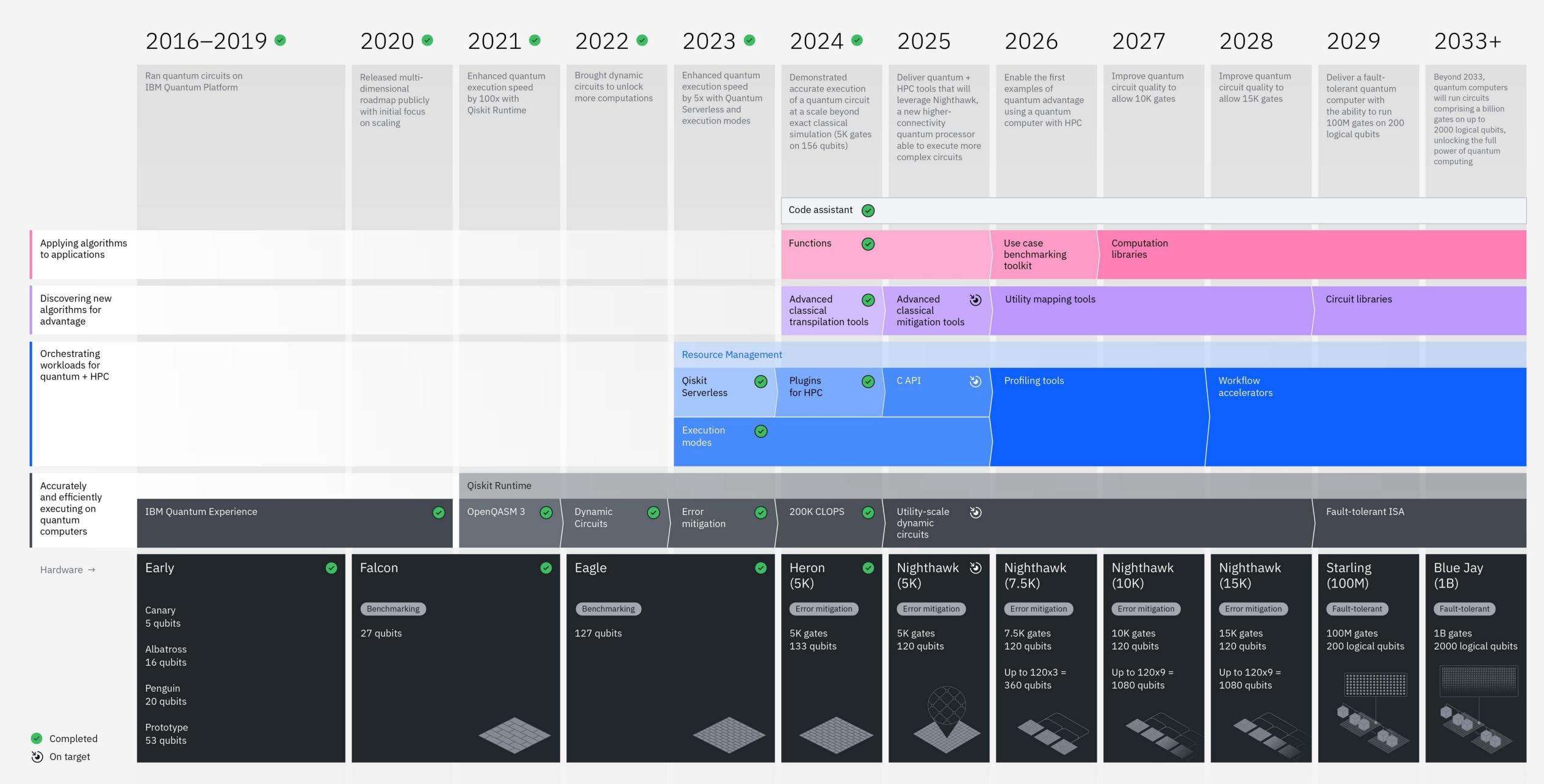
2

Quantum advantage is not something that will happen at a singular moment in time.

It is a hypothesis that is subject to falsification.

There are two important nuances to this →

Development Roadmap



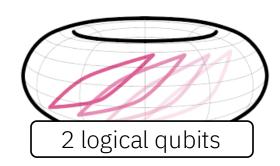
#### IBM fault-tolerant quantum computing roadmap

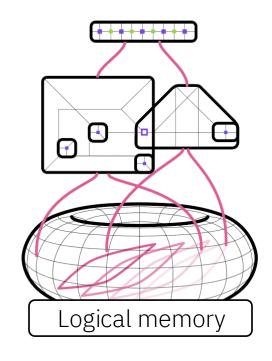
Loon (2025)

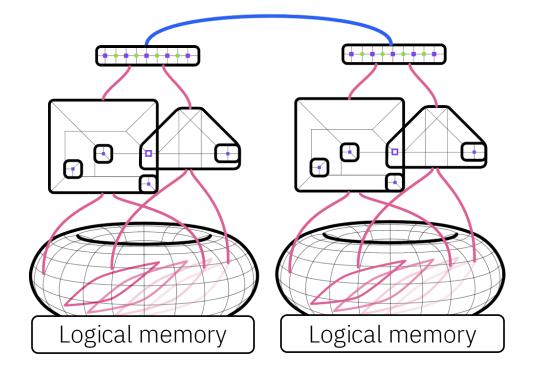
Kookaburra (2026)

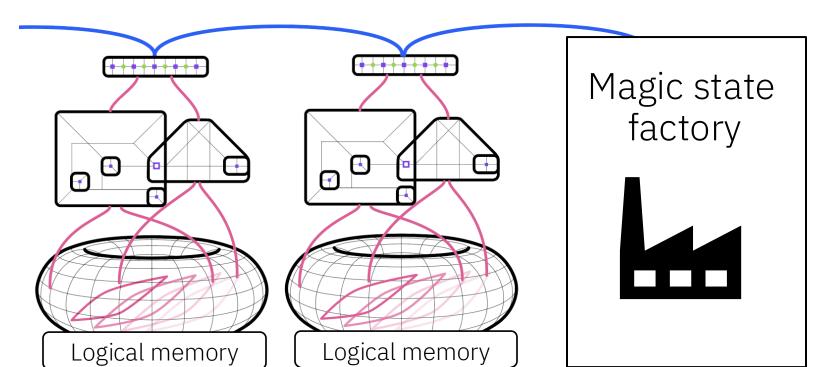
Cockatoo (2027)











6-way connectivity

c-coupler demo

Automated design

"Long" c-couplers

Real-time decoding

LPU (+~100 qubits) for logical operations Two blocks of gross code + LPU

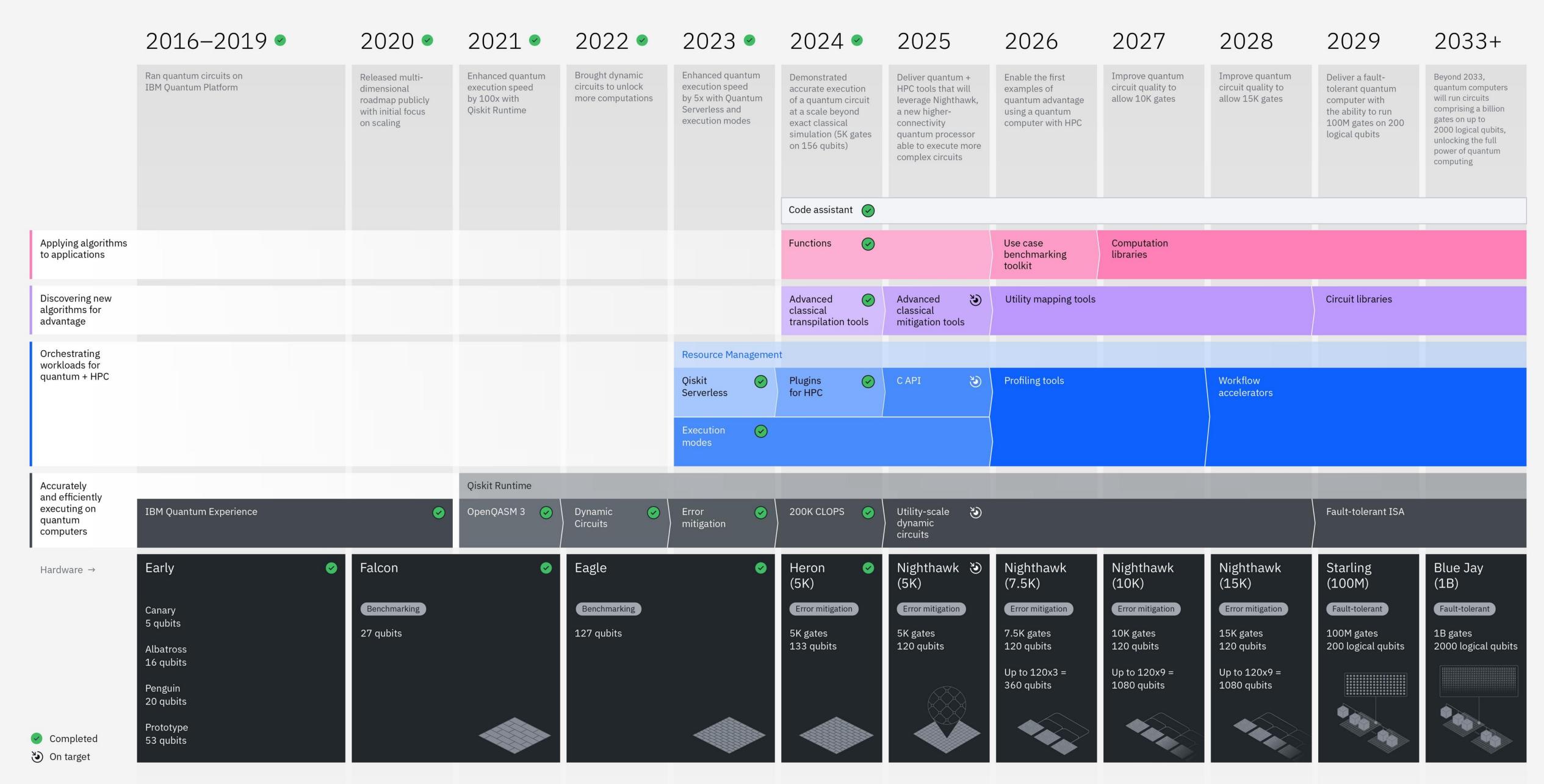
Module-to-module logical communication over l-couplers

Multiple blocks of gross code + LPUs

Universal computation with magic state distillation

Logical memory

Development Roadmap



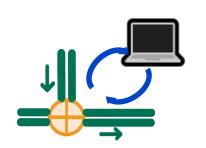
# What are Dynamic Circuits?

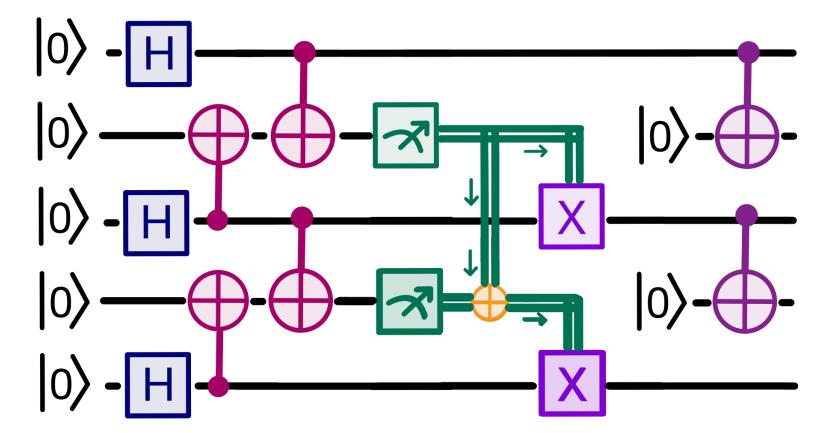
Dynamic circuits are quantum circuits that incorporate

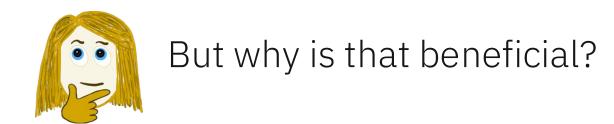
Mid-circuit measurements



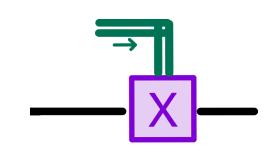
Classical calculations based on the measurement results







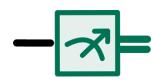
Feed-forward operations



# Why Dynamic Circuits?

Dynamic circuits are quantum circuits that incorporate

# Mid-circuit measurements



Conditional reset:

$$|\psi\rangle - |x| = |x|$$

$$= |\psi\rangle$$
— $|0\rangle$ —

# Classical calculations based on the measurement results

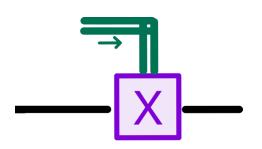


- For certain problems, quantum computers are exptected to give an advantage
- But: in general, classical computations are much faster, have much higher fidelity and no connectivity constraints
  - → Hybrid classical-quantum algorithms like VQE etc



But measurements make the state collapse?

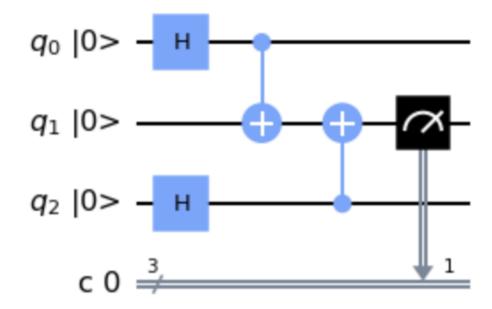
# Feed-forward operations



- Famous application: Error
   Correction
  - → Check for errors and directly correct them
- Use classical calculations as information transfer to spread correlations faster
  - → Create long-range entanglement in a shallow quantum circuit

# Dynamic Circuits for Long-Range Entanglement

Let us consider the following example:

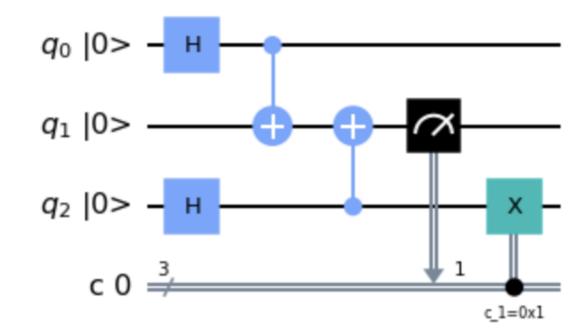


post-measurement state either

$$\frac{1}{\sqrt{2}}\left(|00\rangle_{AC}+|11\rangle_{AC}\right) \text{ or } \frac{1}{\sqrt{2}}\left(|01\rangle_{AC}+|10\rangle_{AC}\right)$$

actually a mixed state

• If we apply a conditional gate:

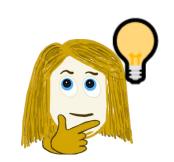


Final state

$$\frac{1}{\sqrt{2}}\left(|00\rangle_{AC} + |11\rangle_{AC}\right)$$

- > using dynamic circuits we can create entanglement without a direct link in constant depth
- especially useful for quantum devices with limited connectivity!

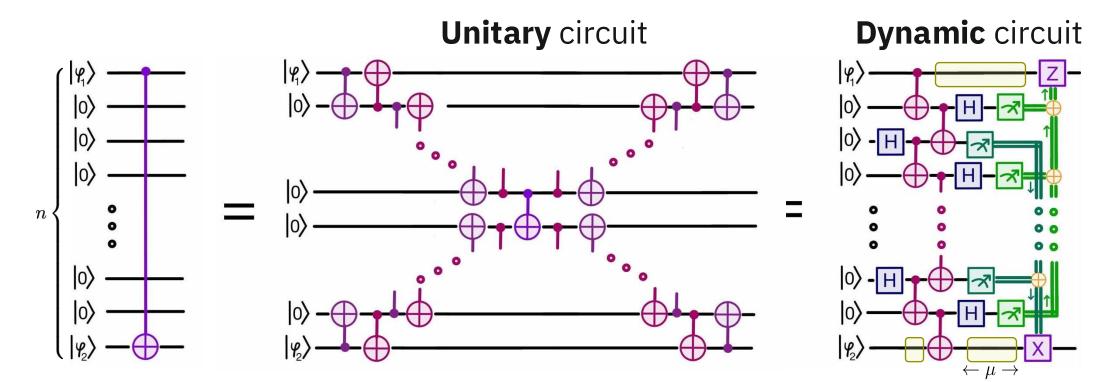
But why real-time feed-forward and not just post-processing?



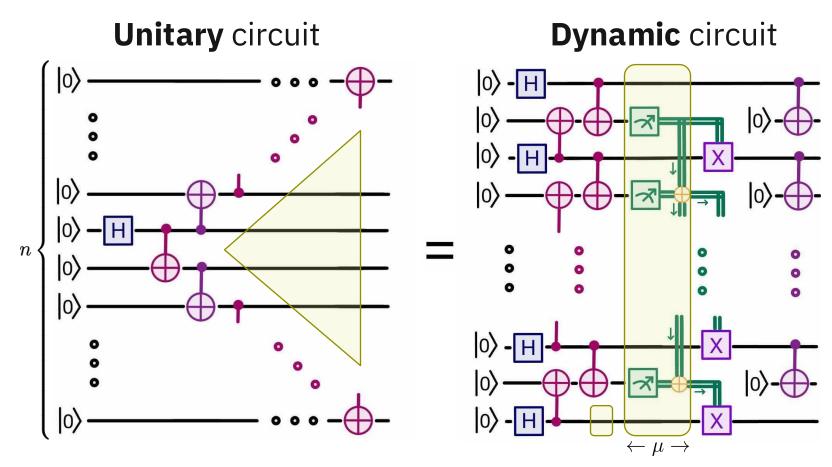
Post-processing only works if we can simulate the different outcomes through the whole circuit!

# Dynamic Circuits for Long-Range Entanglement

- Implementation of long-range entangling gates due to local connectivity
  - to entangle two qubits that are distance n away we need depth  $\mathcal{O}(n)$



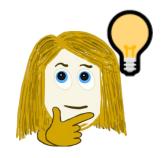
- Preparation of long-ranged entangled states
  - > to entangle n qubits, we need at least depth  $\mathcal{O}(\log n)$



EB, V. Tripathi, D. Wang, P. Rall, E. Chen, S. Majumder, A. Seif, Z. Minev, https://arxiv.org/abs/2308.13065

> Using dynamic circuits, we can create long-range entanglement in a constant depth quantum circuit

But entanglement cannot spread faster than the information light-cone?

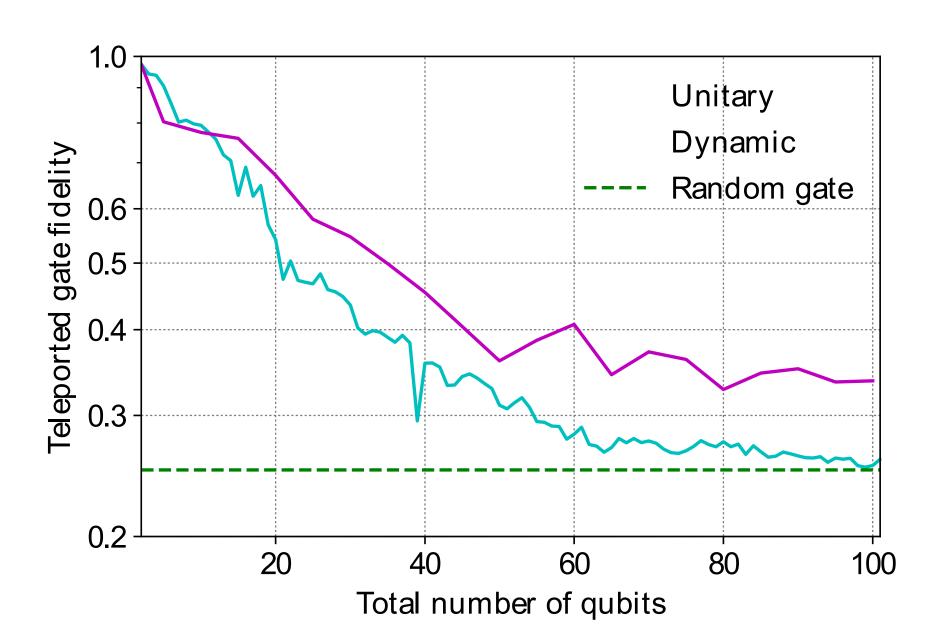


True, but the information transfer in this case is performed by the classical calculation!

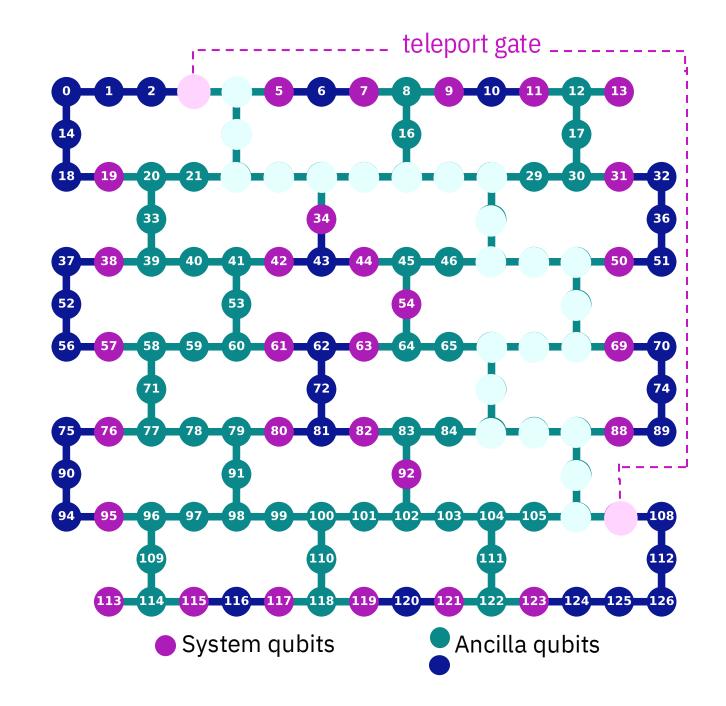
### Experimental Results

Joint work with Vinay Tripathi, Derek Wang, Patrick Rall, Edward Chen, Swarnadeep Majumder, Alireza Seif & Zlatko Minev

 CNOT gates over large distances are more efficiently executed with dynamic circuits than unitary ones



• Outlook: Teleporting gates can provide a workaround for effective all-to-all connectivity

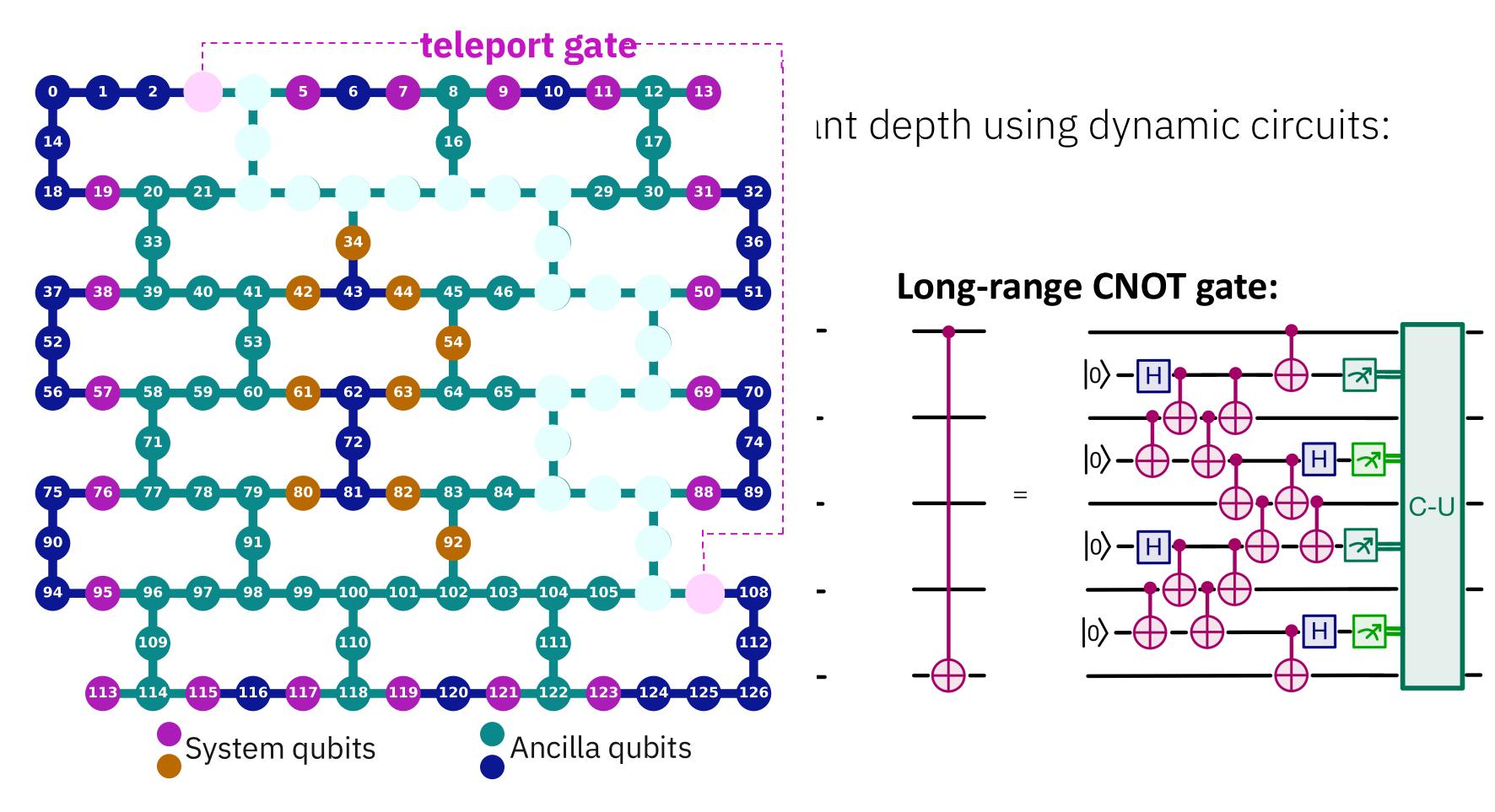


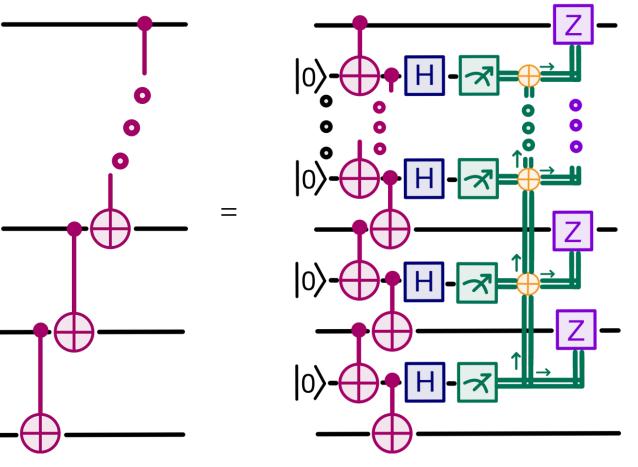
Dynamic circuits are a promising feature to help overcome current limitations of hardware, especially for increasing number of qubits

# Measurement-Based Long-Range Entangling Gates in Constant Depth

Joint work with Stefan Wörner

- Setting: n+1 system qubits
- > Constructions for quantum s



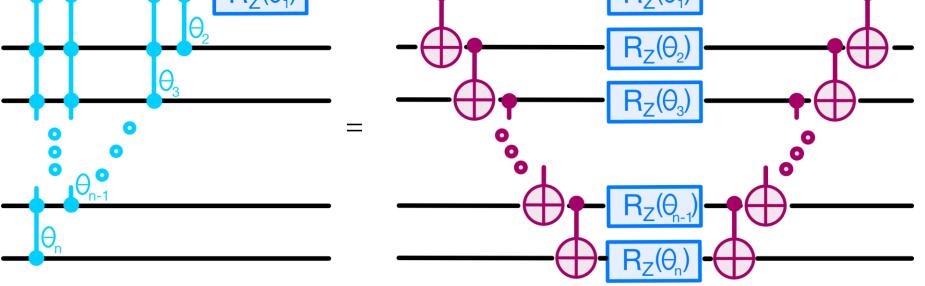


# Measurement-Based Long-Range Entangling Gates in Constant Depth

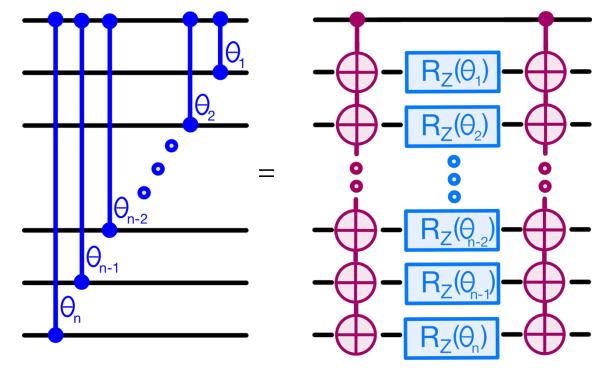
Joint work with Stefan Wörner

> These constructions can be combined to construct multi-qubit (parametrized) rotation gates:

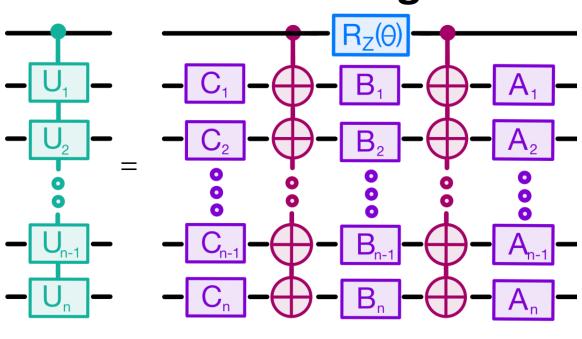
# Multi-qubit $R_{ZZ..Z}$ rotations: $R_{Z}(\theta_{1}) = R_{Z}(\theta_{2})$ $R_{Z}(\theta_{2})$



#### Parallel R<sub>ZZ</sub> rotations:



#### **Generalized fan-out gate:**



## Experimental Results

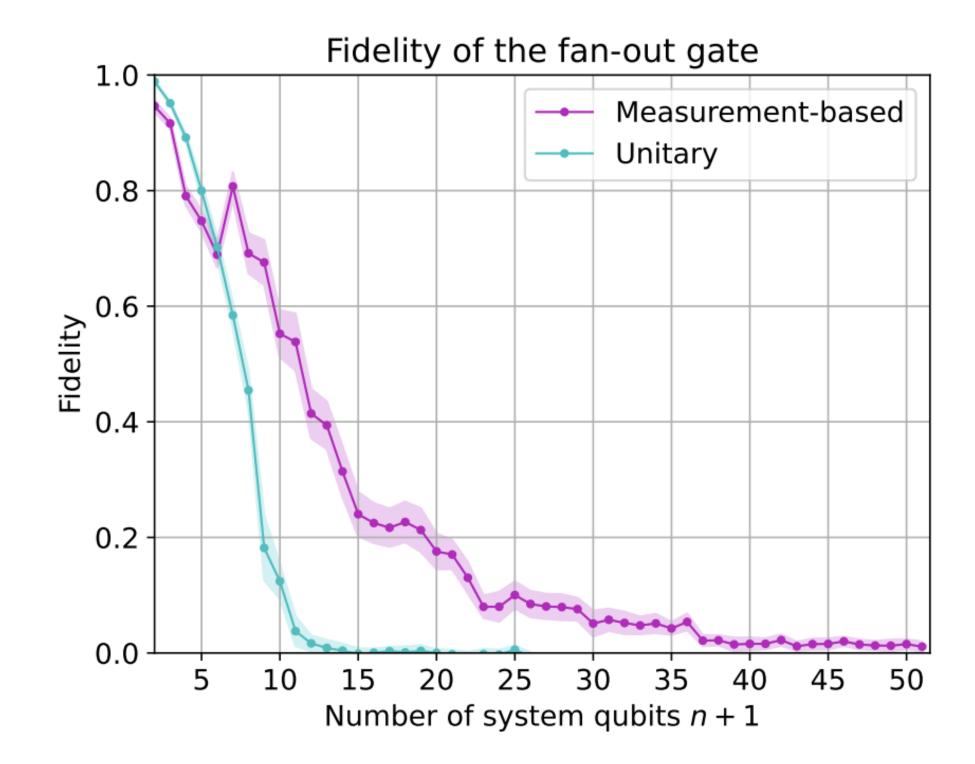


FIG. 5. Experimental results on  $ibm_kyiv$  [13] implementing the fan-out gate on n+1 qubits (2n+1) total qubits in the measurement-based implementation).

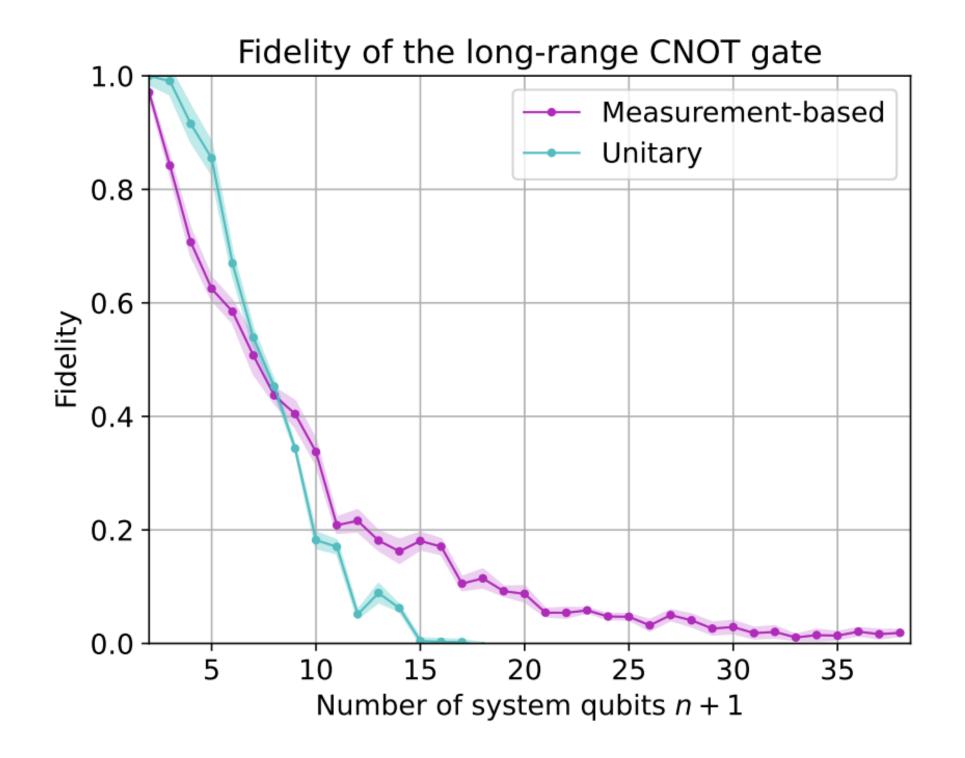


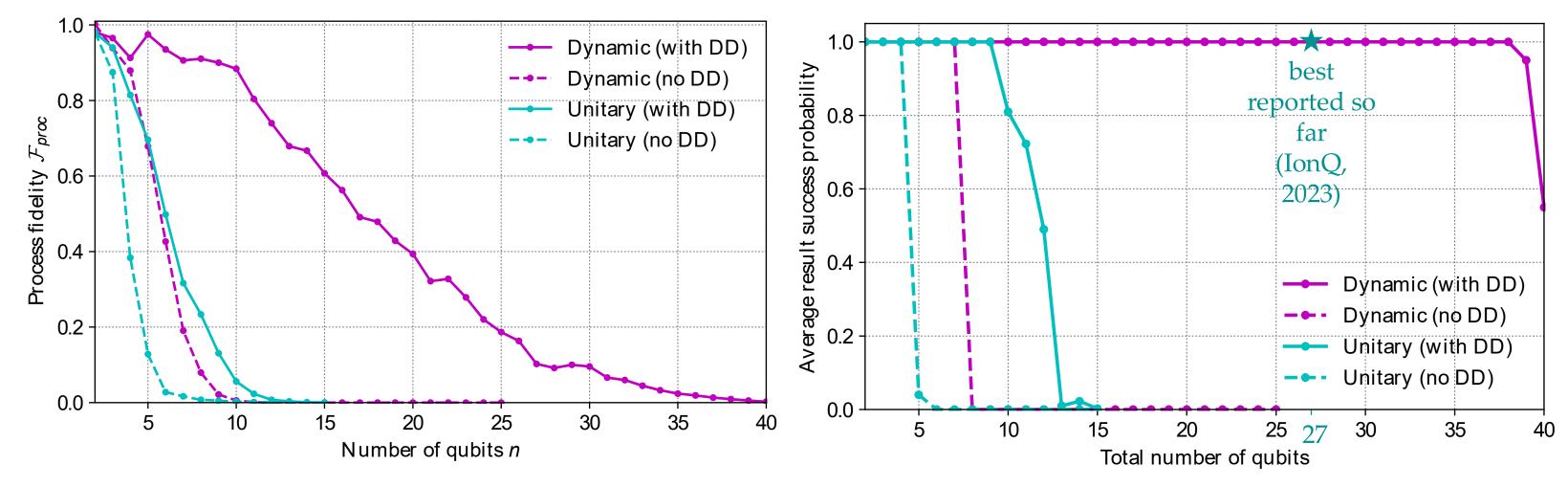
FIG. 6. Experimental results on  $ibm_kyiv$  [13] implementing the long-range CNOT gate on n+1 qubits (2n+1) total qubits in the measurement-based implementation).

> The measurement-based protocol outperforms the unitary one for larger number of qubits

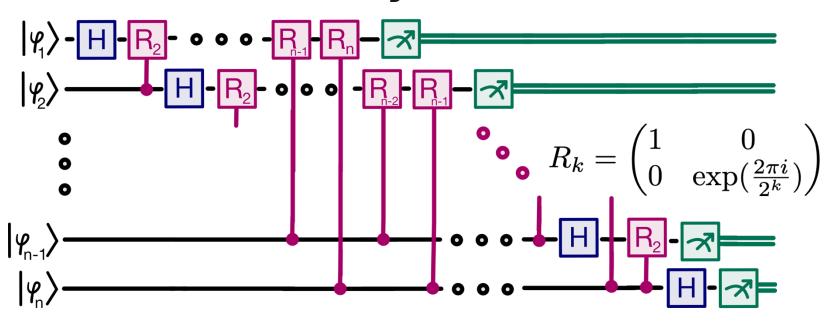
# Quantum Fourier Transform using Dynamic Circuits

Joint work with Vinay Tripathi, Alireza Seif, Daniel Lidar & Derek Wang

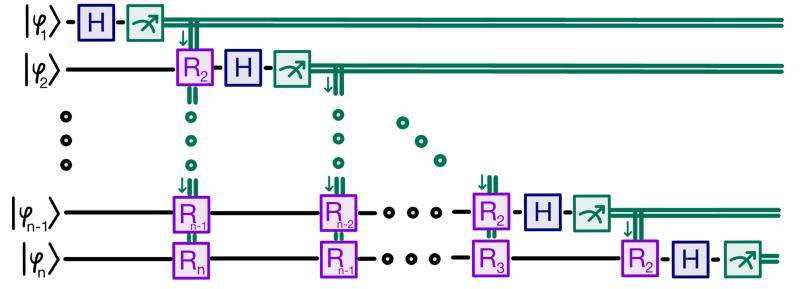
- Important subroutine of many quantum algorithms, e.g. Shor's Algorithm or Quantum Phase Estimation. It is often followed by a measurement and can then be drastically simplified using dynamic circuits.
- Instead of  $\mathcal{O}(n^2)$  two-qubit gates in the standard unitary circuit, we only require  $\mathcal{O}(n)$  mid-circuit measurements in the dynamic circuit without any connectivity constraints.



#### **Unitary** circuit

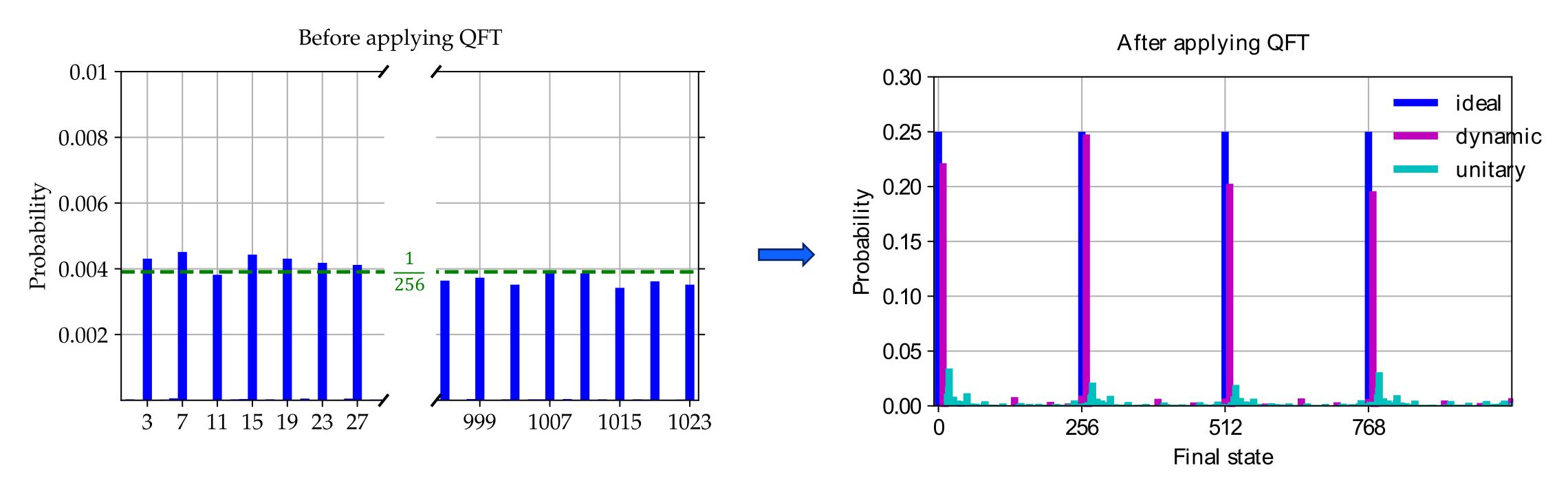


#### **Dynamic** circuit



## Experimental Results

• To visualize the improvement of QFT+M with dynamic circuits, we prepare a simple periodic state on 10 qubits and compare the ideal, unitary and dynamic circuit implementation of QFT+M applied to that state



> Implementation with dynamic circuits resembles ideal case, while unitary circuits result in a flatter distribution

# Approximate Quantum Fourier Transform in Logarithmic Depth on a Line

Joint work with David Sutter & Stefan Wörner

Approximate the Quantum Fourier Transform with error  $\varepsilon$  in depth  $\mathcal{O}\left(\log\frac{n}{\varepsilon^2}\right)$  using only  $\mathcal{O}(n)$  qubits on a line!

#### Construction:

For most inputs, we require only two operations:

i. Quantum Fourier state computation (QFS):

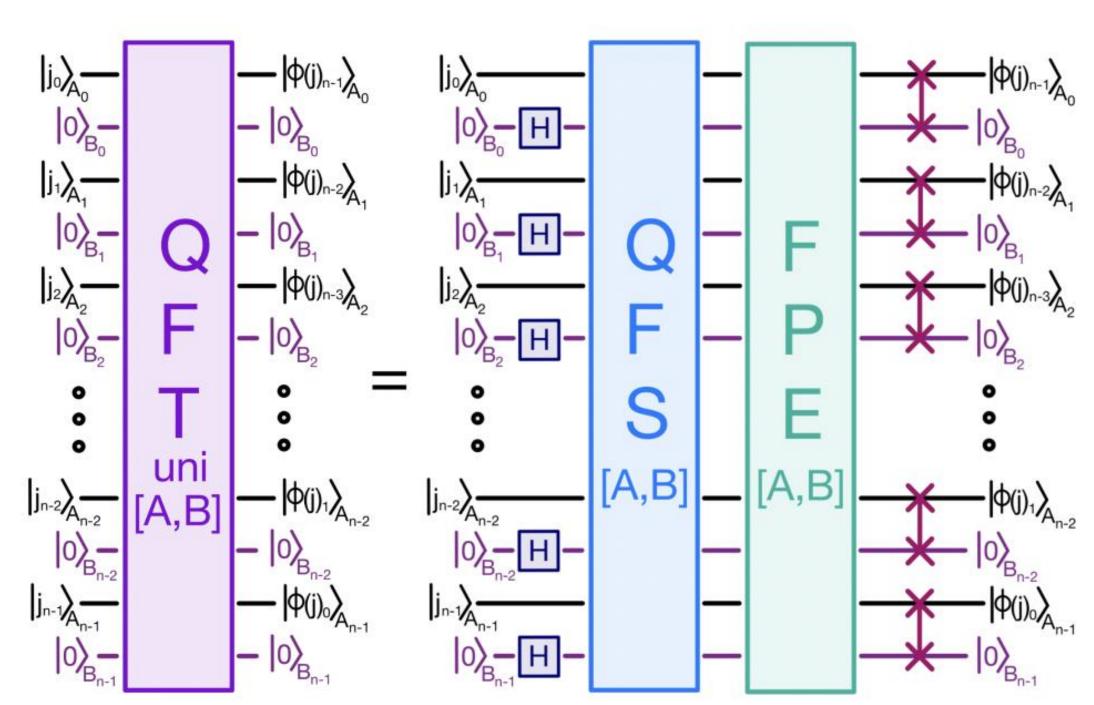
$$|j\rangle|\phi(b)\rangle \rightarrow |j\rangle|\phi(b+j)\rangle$$

ii. Fourier phase estimation (FPE):

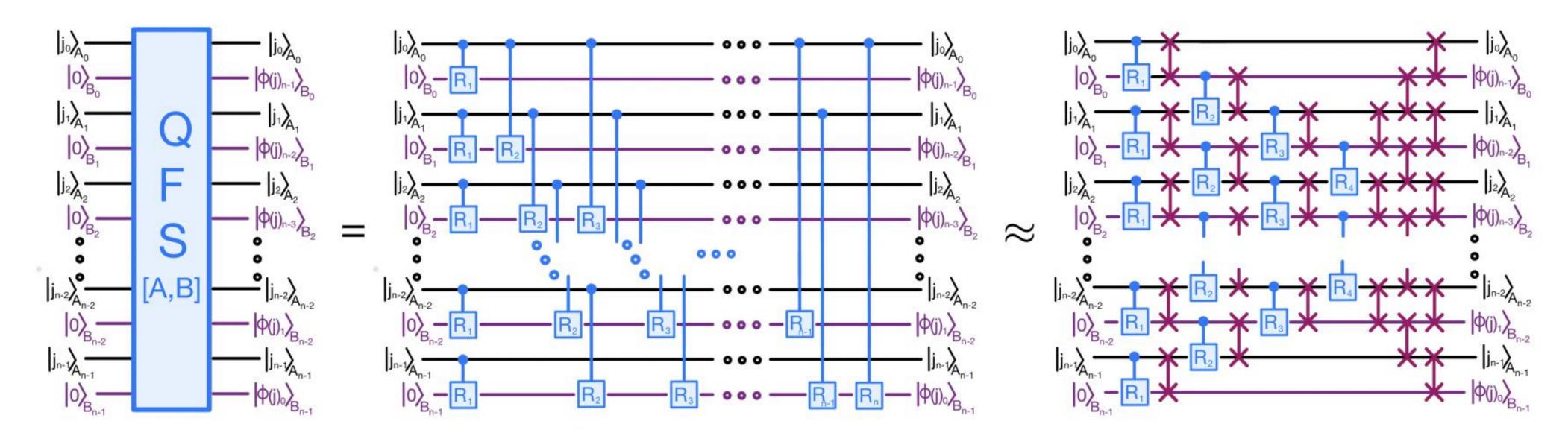
$$|b\rangle|\phi(j)\rangle \rightarrow |b\oplus j\rangle|\phi(j)\rangle$$

together with a Hadamard transform and SWAP gates:

$$|j\rangle|0\rangle \xrightarrow{H} |j\rangle|\phi(0)\rangle \xrightarrow{\text{QFS}} |j\rangle|\phi(j)\rangle \xrightarrow{\text{FPE}} |0\rangle|\phi(j)\rangle \xrightarrow{\text{SWAP}} |\phi(j)\rangle|0\rangle$$



# Quantum Fourier state computation (QFS)



- Exact QFS:  $|j\rangle_A|\phi(b)\rangle_B \overset{\mathrm{QFS}_{AB}}{\to} |j\rangle_A|\phi(b+j)\rangle_B$
- Approximate by neglecting small rotations, i.e. all phase gates  $R_k:=egin{pmatrix} 1 & 0 \ 0 & \mathrm{e}^{2\pi\mathrm{i}/2^k} \end{pmatrix}$  for  $k>k_{\mathrm{max}}=O(\log rac{n}{arepsilon})$
- $ightharpoonup 
  m A \ unitary \ QFS^{(arepsilon)} \ \ acting on \ \ \mathcal{S} := \left\{ |\psi\rangle_{ABE} \in \mathcal{H}_{ABE} : |\psi\rangle_{ABE} = \sum_{j=0}^{N-1} \sum_{m=0}^{M-1} \alpha_{j,m} |j\rangle_{A} |0\rangle_{B} |m\rangle_{E} \right\}$  with  $dist_{\mathcal{S}}(QFS_{AB}, QFS_{AB}^{(arepsilon)}) \le \varepsilon$  can be implemented on a line of 2n qubits with nearest-neighbor

connectivity with depth  $O(\log \frac{n}{\varepsilon})$ 

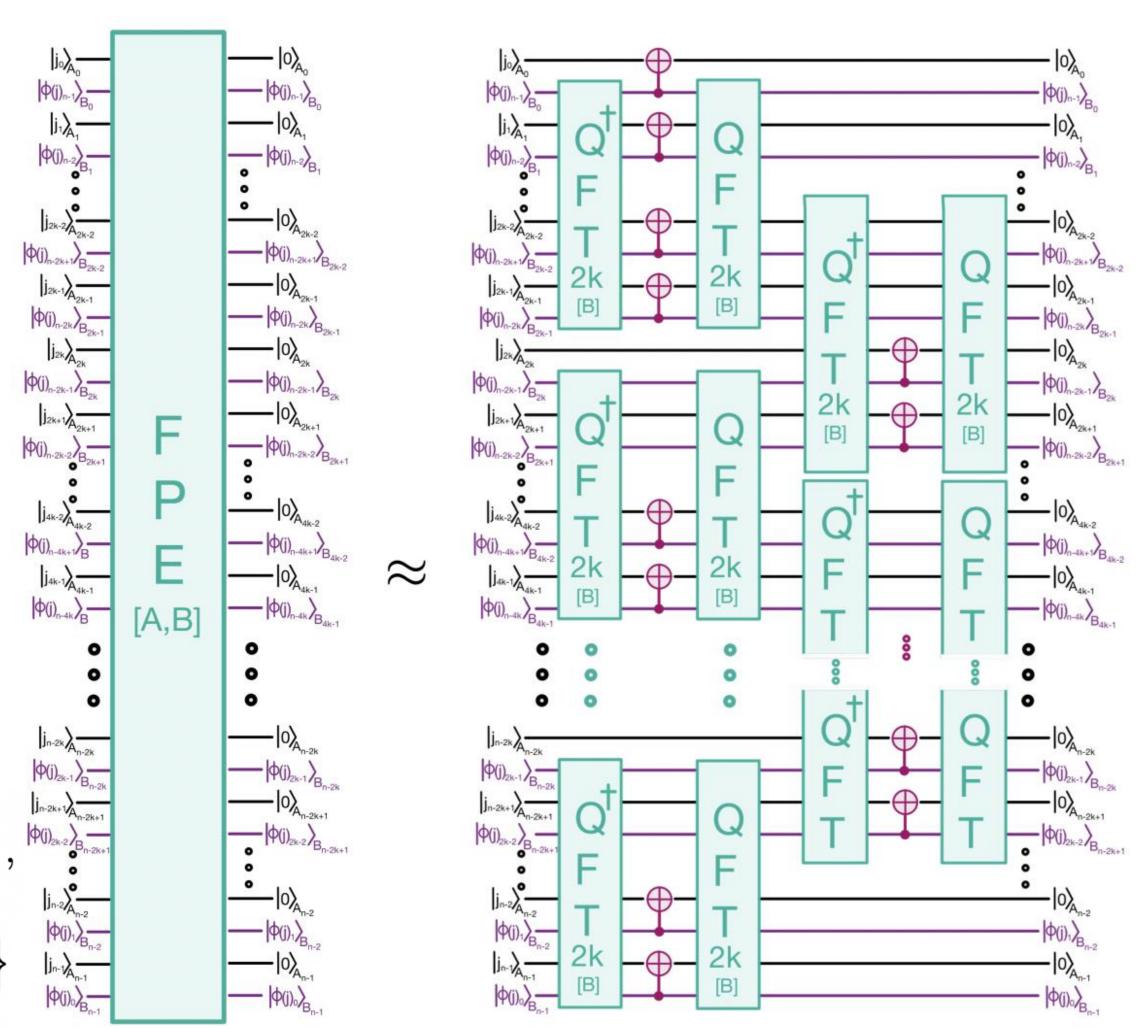
# Fourier phase estimation (FPE)

- Exact FPE:  $|b\rangle_A|\phi(j)\rangle_B \overset{\mathrm{FPE}_{AB}}{\longrightarrow} |b\oplus j\rangle_A|\phi(j)\rangle_B$
- Approximation: estimate  $|j\rangle$  by small, but exact QFTs that are applied in parallel on 2k qubits each, where  $k=\mathcal{O}(\log n)$
- A unitary  $\mathrm{FPE}^{(\varepsilon)}$  with  $\mathrm{dist}_{\mathcal{T}^{(p,q)}_{\mathrm{uni}}}(\mathrm{FPE}_{AB},\mathrm{FPE}^{(\varepsilon)}_{AB}) \leq \varepsilon$  can be implemented on a line of 2n qubits with nearestneighbor connectivity with depth  $O(\log \frac{n}{\varepsilon^2})$

$$! \quad \text{BUT:} \quad \mathcal{T}_{\text{uni}}^{(p)} := \left\{ |\psi\rangle_{ABE} \in \mathcal{H}_{ABE} : |\psi\rangle_{ABE} = \sum_{j=0}^{N-1} \sum_{m=0}^{M-1} \beta_{j,m} |j\rangle_{A} |\phi(j)\rangle_{B} |m\rangle_{E} \right.,$$

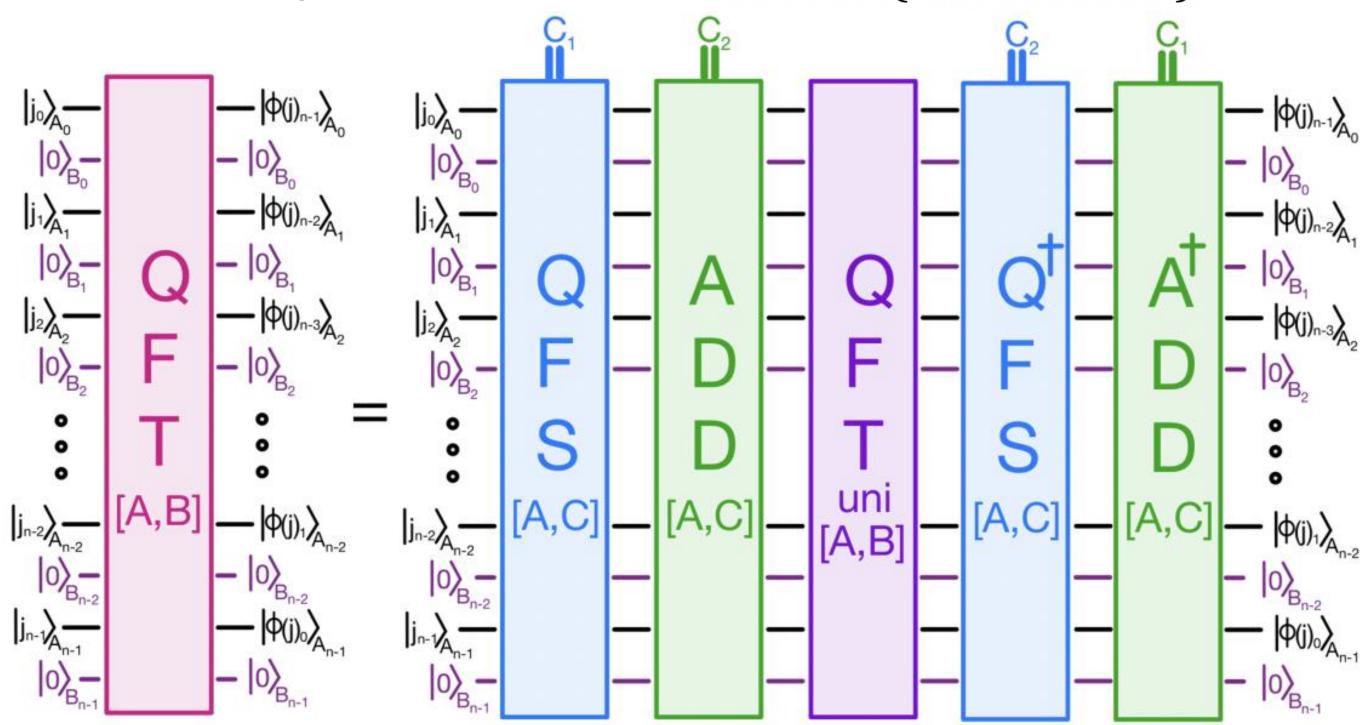
$$\left| \sum_{m} \beta_{j,m} \beta_{\ell,m}^{*} \right| \leq \frac{p(n)}{N} \delta_{j,\ell} \, \forall j, \ell \left. \right\}$$

> need somewhat uniform inputs



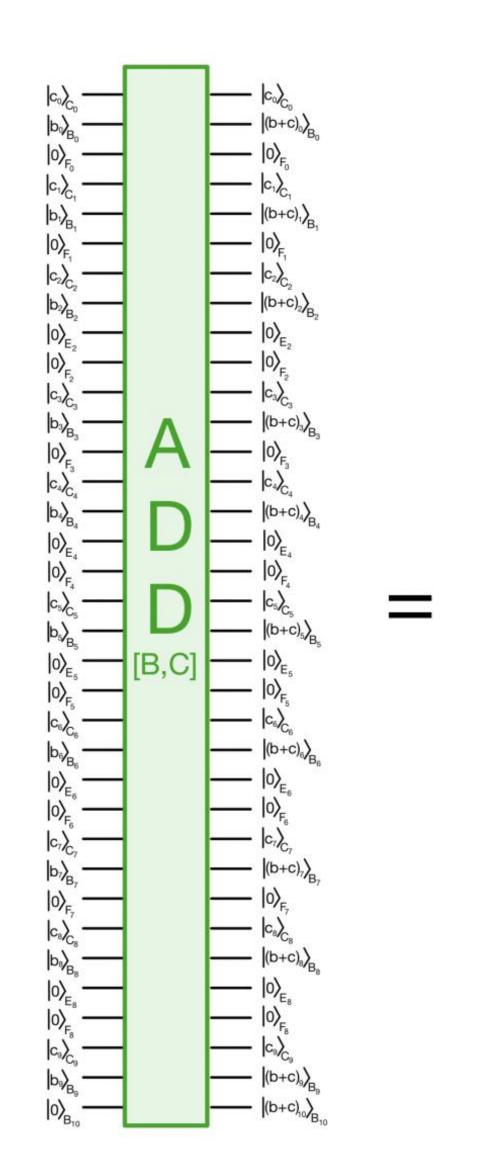
# Constructing uniform inputs

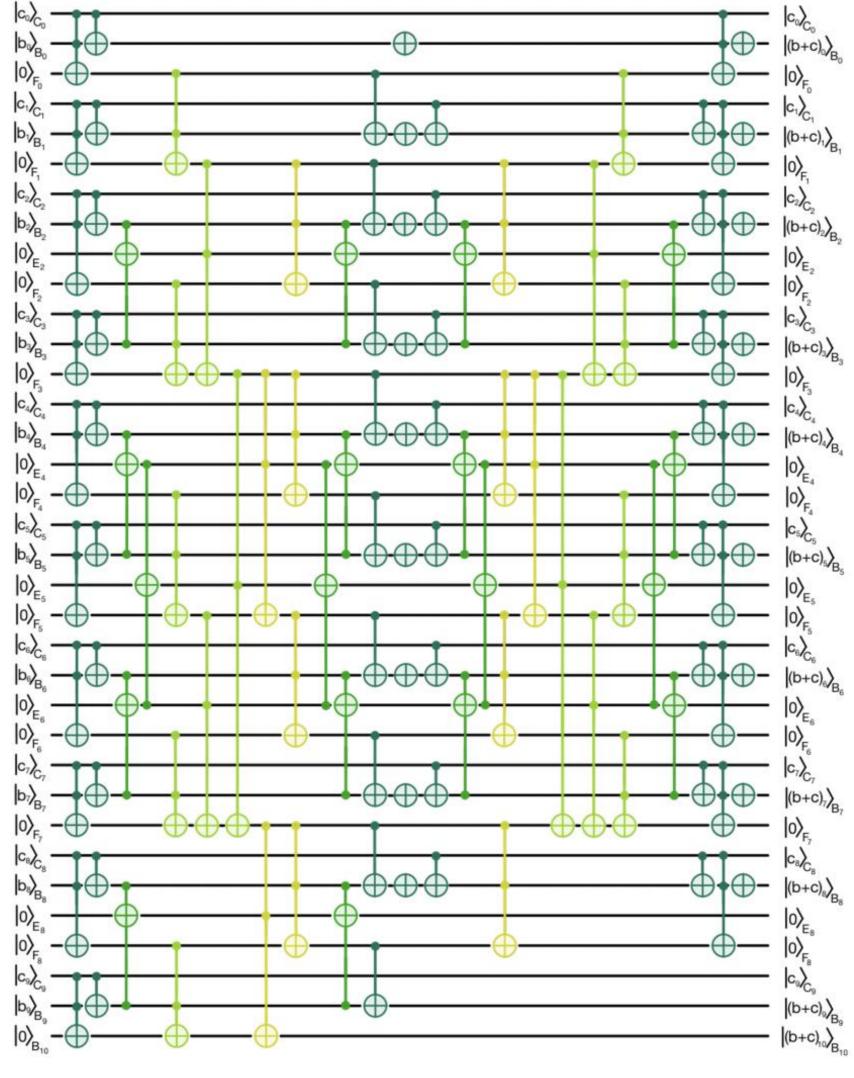
- Problem: approximate FPE works for most inputs, but not all
- > Constructing a somewhat uniform input is guaranteed to work
- ightharpoonup Requires another operation: Quantum adder (ADD):  $|b\rangle|c\rangle \to |b+c\rangle|c\rangle$  as well as a classical register C initialized with two randomly chosen numbers  $c_1,c_2\in\{0,\ldots,N-1\}$
- ightharpoonup Constructed an adder that can be implemented on a line of  $\mathcal{O}(n)$  qubits with nearest-neighbor connectivity in logarithmic depth using dynamic circuits



# Quantum Adder (ADD)

- ADD:  $|b\rangle|c\rangle \stackrel{\mathrm{ADD}}{\to} |b+c\rangle|c\rangle$  also equivalent to  $|\phi(j)\rangle|\phi(\ell+j)\rangle \to |\phi(j)\rangle|\phi(\ell)\rangle$
- Draper et al.'s logarithmic-depth quantum carrylookahead adder seems to require all-to-all connectivity
- Notice that in each of the  $\mathcal{O}(\log n)$  layers, the long-range gates are not overlapping and thus we can use dynamic circuits to implement them in parallel even on a 1D line
- The ADD mapping can be implemented on a line of 5n qubits with nearest-neighbor connectivity with depth  $\mathcal{O}(\log n)$





T. G. Draper, S. A. Kutin, E. M. Rains, and K. M. Svore, A logarithmic-depth quantum carry-lookahead adder

# Improved Implementation

- AQFT for uniform inputs: A unitary  $\operatorname{QFT}_{\mathrm{uni}}^{(\varepsilon)}$  acting on uniform states with  $\operatorname{dist}_{\mathcal{S}_{\mathrm{uni}}^{(p)}}(\operatorname{QFT},\operatorname{QFT}_{\mathrm{uni}}^{(\varepsilon)}) \leq \varepsilon$  can be implemented on a line of 2n qubits with nearest-neighbor connectivity with depth  $O(\log \frac{n}{\varepsilon^2})$ 
  - > Guaranteed to work for uniform inputs, but also works for most random input states
  - > May be very relevant in practice, as it can always be the first try if the result is verifiable
- AQFT for general inputs: A unitary  $\operatorname{QFT}^{(\varepsilon)}$  vith  $\operatorname{dist}_{\mathcal{S}}(\operatorname{QFT}_{AB},\operatorname{QFT}_{AB}^{(\varepsilon)}) \leq \varepsilon$  can be implemented on a line of 4n qubits with nearest-neighbor connectivity with depth  $O(\log \frac{n}{\varepsilon^2})$
- Use the principle of deferred measurement on the (ideally) uncorrelated ancilla qubits for uncomputation
- > Running QFT "backwards" allows to implement the QFS exactly in constant depth
- Potential optimization: recursively replace small QFTs in the FPE by this approximate QFT to reduce depth of FPE to  $\mathcal{O}(\log \log n)$   $\rightarrow$  analysis of errors and trade-offs due to larger pre-factors required

### Conclusion & Outlook

- Dynamic circuits can help overcome connectivity constraints, reduce the depth and thereby drastically improve the compilation of quantum algorithms
- Their effect will become more significant as we scale up the number of qubits
- Already now we can see the benefit of leveraging dynamic circuits in optimizing the compilation of some quantum algorithms, many more applications to explore!
  - > Improved compilation to overcome current limitations of hardware
  - > Optimizing the resources required in the fault-tolerant compilation

Thank you for your attention!